



Water Resource System Optimization by Geometric Programming

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RESEARCH PROJECT TECHNICAL COMPLETION REPORT

Project Number A-011-TEX

September 1969 - August 1971

Agreement Number

14-01-0001-1864

14-31-0001-3044

14-31-0001-3244

WATER RESOURCE SYSTEM OPTIMIZATION

BY GEOMETRIC PROGRAMMING

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The work upon which this publication is based was supported in part by funds provided by the United States Department of the Interior, Office of Water Resources Research, as authorized under the Water Resources Research Act of 1964, P.L. 88-379.

Technical Report No. 34
Water Resources Institute
Texas A&M University

February 1971

PREFACE

The research described in this report was supported by the U. S. Department of the Interior, Office of Water Resources Research and the Texas A&M University Water Resources Institute. The aim of this research was to study the applicability and usefulness of geometric programming in water resource planning and design activities. In order to study the use of geometric programming in practical problem solving, a computer algorithm was developed. A variety of water resources problems were solved using the algorithm.

The authors believe that there is a significant need for techniques which will permit the optimum design and operation of either proposed or existing water filtration systems. It was also believed that, because the cost models and constraints for water filtration design problems are generalized polynomials, geometric programming would be a good optimization technique to use. Because water filtration problems have not been the subject of optimization analyses in the past to any significant degree, some time was spent in developing cost models and formulating the problem. This effort and the optimization analysis combining linear and geometric programming was of such interest that it is described rather extensively in the report. The report includes a tutorial explanation of geometric programming, a description of the computer algorithm used, the description of the water filtration cost model, and an explanation of the use of geometric programming in an application.

Appreciation is expressed to the numerous colleagues who assisted in the satisfactory completion of this research. Thanks are due to Mr. Thomas H. Edwards for his help during the initial phases of the research. The assistance and support of Dr. J. R. Runkles and the staff of the Water Resources Institute at Texas A&M University is gratefully acknowledged. Special thanks are due to Mrs. Jane Teske who spent long hours in typing the manuscript.

TABLE OF CONTENTS

	Page
PREFACE	ii
TABLE OF CONTENTS	iv
LIST OF TABLES.	vi
LIST OF FIGURES	vii
Chapter	
I INTRODUCTION	1
Optimization in Water Resources Analyses	1
Mathematical Programming	2
Development of Geometric Programming	3
II GEOMETRIC PROGRAMMING.	6
Posynomial Functions	6
Optimizing Posynomials	7
Constrained Optima	12
Degrees of Difficulty.	13
Computer Algorithm	16
Generalized Polynomials	17
III COMPUTER ALGORITHM	22
Generalized Notation	22
Main Program	23
Subroutine MAT	25
Subroutine LETTO	25
Subroutine VAL	29
Other Subroutines.	29
Input.	30
Output	32
Error Messages	33
IV WATER FILTRATION THEORY.	34
The Filtering Process.	34
Physical Removal Mechanisms.	34
Chemical Removal Mechanisms.	37
Empirical Studies.	38
Multimedia Filters	40

Chapter		Page
V	FILTER COST MODEL.	45
	Reducing the Problem	45
	Objective Function	47
	Constraints.	49
	An Optimum Solution.	56
VI	A SAMPLE PROBLEM	57
	Constant Values.	57
	Model for Sample Problem	59
	Geometric Programming Solution	61
	Linear Programming Solution.	62
	Summary of Results	63
VII	CONCLUSIONS AND RECOMMENDATIONS.	66
	Filter Cost Model.	66
	Geometric Programming Optimization	66
	Combining Optimization Techniques.	67
	Future Extensions.	68
	LIST OF REFERENCES	69
	APPENDIX	
	A Program Listing	74
	B Example of Output	108

LIST OF TABLES

Table	Page
6.1 OUTPUT COMPARISONS	65

LIST OF FIGURES

Figure		Page
3.1	PROGRAM GPS	24
3.2	SUBROUTINE MAT FLOW DIAGRAM	26
4.1	PERFORMANCE CURVE 1 FOR A SYNTHETIC MUDDY WATER	41
4.2	PERFORMANCE CURVE 2 FOR A SYNTHETIC MUDDY WATER	42
4.3	COMPARISON OF FILTER MEDIA.	43
5.1	MULTIMEDIA FILTER	46

CHAPTER I

INTRODUCTION

Optimization in Water Resources Analyses

Water resources planners and systems analysts are continually confronted with many complex optimization problems. Two major factors contribute to this problem. Firstly, mathematical modeling and system description capabilities in water resources have increased considerably. Secondly, planning and design analyses have become more complicated because of a desire by planners to represent more completely the problem to be solved.

In a recent appraisal of Federal Water Resources Research activities, research concerned with water resource planning methodology was labeled as the "most promising area of research" [22]. In an effort to meet the Congressional decree [54] for "optimum" water resources planning, the literature is filled with illustrations of the applicability and usefulness of optimization methods in water resources analyses [38]. Linear programming [53], dynamic programming [41,40], and nonlinear programming [19] have been shown to be useful in water resource optimization problems. These techniques have been and are being coupled with mathematical modeling and system simulation to produce practically useful planning algorithms.

The purpose of this report is to describe a new optimization technique which can be extremely useful in solving water resources optimization problems. This new and potentially powerful technique

is called *geometric programming*. It is one of a class of mathematical programming techniques. Mathematical programming discussed extensively elsewhere [56,23] refers to a class of optimization problems dealing with the allocation of limited resources. Problems are described with an objective function which measures system effectiveness and constraints which describe limitations on the solution. If either the objective function or constraints are nonlinear, a nonlinear programming problem results.

A variety of algorithms are available for solving nonlinear programming problems. However, when the constraints are nonlinear, the objective function is higher than second degree or the problem of high dimensionality; most of the available techniques are not computationally efficient. Little previous research has been aimed at making geometric programming of use to water resource analysts. The purpose of this research was to explore ways of making geometric programming more useful to water resource systems analysts. Specifically, this research has as its objectives the development of a computer algorithm for solving geometric programming problems and the specific application of this algorithm in the solution of a water renovation and treatment process.

Mathematical Programming

A general statement of the mathematical programming problem is

$$\begin{aligned} &\text{maximize/minimize: } y = g(x) \\ &\text{subject to: } g_i(x) \begin{matrix} \geq \\ < \end{matrix} b_i \quad i=1, \dots, N, \\ &x \geq 0. \end{aligned} \quad (1.1)$$

A maximization problem is equivalent to the minimization of the negative of the objective function shown above. The forms of $g(x)$ are important in determining the method used to solve the problem. For example if the $g(x)$ are all linear functions of non-negative variables, the problem is called a linear programming problem, and the method of solution is the simplex method [56].

The general theory of mathematical programming was developed by Kuhn and Tucker [35]. The Kuhn-Tucker theory applies both the linear and nonlinear programming problems. The theory utilizes a generalization of the LaGrangian method of undetermined multipliers. It extends the LaGrangian method to inequality constraints. Kuhn-Tucker theory relies on concepts of convexity and duality.

In general, according to the theory, if a constrained objective function (primal program) is convex, then, for all $x > 0$, there exists an equivalent dual program, involving a dual function with a maximum at the primal minimum point. This theory of duality provides an indirect means of arriving at the optimum, since the dual program may be easier to solve than the primal and provides an easy way to check whether or not the optimum has indeed been reached.

Development of Geometric Programming

Many mathematical models involve variables raised to arbitrary powers. A great many theoretical laws and empirical results are expressed in such terms and others can be reduced to that form. It was the problem of minimizing the sum of component costs equal to design variables raised to arbitrary powers that led to the observation

by Zener [58] that such functions could be minimized almost by inspection, when the number of terms was one greater than the number of variables. These functions of the form shown in Equation 1.2 can be

$$y = \sum_{t=1}^T c_t \prod_{n=1}^N x_n^{a_{tn}} \quad (1.2)$$

classed as *generalized polynomials*. However, Zener and Duffin [60] established a theory in which the a_{tn} and x_n above could be any real number but the c_t were restricted to be positive. They called these functions *posynomials*. Following an approach akin to dimensional analysis, weights were chosen to make the product of component costs non-dimensional in the dual function. This product, Zener observed, was in fact the minimum cost.

Consultation with Duffin extended the theory to posynomials of an arbitrary number of terms [17], and related it [16] to Cauchy's inequality [24]. This theoretical connection to Cauchy's inequality, also called the *arithmetic-geometric mean inequality*, led to the name *geometric programming*.

Negative coefficients were forbidden the functions, and the sense of inequality was restricted, by the convexity requirement. Extension of geometric programming to negative coefficients and arbitrary inequalities demanded a release from convexity, accompanied necessarily by a release from the arithmetic-geometric mean inequality. The posynomial theory involved raising terms to fractional powers, thus the terms could not involve negative constants.

This extension was made, utilizing signum functions and Lagrangian theory, for the equality constraint case [55], and then for the inequality constraint case [47]. Unfortunately, this release from convexity also released the theory from simple convex duality. The primal and dual functions each may have several maxima, minima, or saddle points. It is this theory upon which the algorithm presented in this report is based. The present theory is also summarized in *Foundations of Optimization* [56].

CHAPTER II

GEOMETRIC PROGRAMMING

Posynomial Functions

Geometric programming was conceived as a means of optimizing the cost in engineering design problems in which the cost is a sum of terms expressing component costs. Zener [58] observed that such a sum may be easily minimized when each cost depends on products of the design variables raised to arbitrary but known powers. Cost functions of this type are of the form

$$y = \sum_{t=1}^T c_t \prod_{n=1}^N x_n^{a_{nt}} \quad (2.1)$$

in which, $c_t \geq 0$, $x_n \geq 0$, and a_{nt} may be any real number. If the non-negativity restriction is removed on c_t , Equation (2.1) represents a generalized polynomial. However, for reasons which will become clear below, Zener [19] restricted the $c_t \geq 0$ and named these particular functions *posynomials*.

The posynomial can be transformed by a change of variables

$$e^{z_n} = x_n, \quad n=1, \dots, N$$

(2.2)

when $x_n > 0$

such that Equation (2.1) becomes

$$y = \sum_{t=1}^T c_t e^{\sum_{n=1}^N a_{tn} z_n} \quad (2.3)$$

For posynomials, Equation 2.3 is convex and therefore has a global minimum. If the sign restriction on c_t is lifted, a stationary point can be found, but there is no guarantee that this point is the global minimum. Formal optimization methods, proofs, and examples are provided by Duffin, Peterson, and Zener [19] and Wilde and Beightler [56]. An introductory summary of basic geometric programming concepts using the notation of Beightler, Crisp, and Meier [2] is provided in the remainder of this chapter.

Optimizing Posynomials

To obtain the minimum of a convex function, the first partial derivatives of y are taken with respect to each of the variables x , and the resulting nonlinear equations are set equal to zero (signifying that the first partial derivatives vanish) as shown in Equation 2.4.

$$\frac{\partial y}{\partial x_k} = \frac{1}{x_k^*} \sum_{t=1}^T a_{tk} c_t \prod_{n=1}^N (x_n^*)^{a_{tn}} = 0, \quad k=1, \dots, N \quad (2.4)$$

Following the conventional procedure, weights shown as w_t in Equation 2.5 and representing the ratio of the individual terms in Equation 2.1 evaluated at the optimum x^* to the optimum cost y^* are defined.

$$w_t = \frac{c_t \prod_{n=1}^N (x_n^*)^{a_{tn}}}{y^*}, \quad t=1, \dots, T \quad (2.5)$$

By definition, the weights sum to unity

$$\sum_{t=1}^T w_t = 1 \quad (2.6)$$

Assuming that y^* and x_k^* are non zero, Equations 2.5 and 2.4 give

$$\sum_{t=1}^T a_{tk} w_t = 0, \quad k=1, \dots, N \quad (2.7)$$

Equations 2.6 and 2.7 are linear in w_t and do not depend on the cost coefficient c_t . Since the weights sum to unity

$$y^* = \prod_{t=1}^T (y^*)^{w_t} = \prod_{t=1}^T \frac{c_t \prod_{n=1}^N (x_n^*)^{a_{tn} w_t}}{w_t} \quad (2.8)$$

Inverting the order of the products in the double product term in Equation 2.8 and applying Equation 2.7, one obtains

$$\begin{aligned}
 \prod_{t=1}^T \prod_{n=1}^N (x_n^*)^{a_{tn} w_t} &= \prod_{n=1}^N \prod_{t=1}^T (x_n^*)^{a_{tn} w_t} \\
 &= \prod_{n=1}^N (x_n^*)^{\sum_{t=1}^T a_{tn} w_t} \\
 &= 1
 \end{aligned} \tag{2.9}$$

Therefore Equation 2.8 can be simplified to become

$$y^* = \prod_{t=1}^T \left(\frac{c_t}{w_t^*} \right)^{w_t^*} \tag{2.10}$$

To find y^* , the w_t values are calculated from Equations 2.6 and 2.7 and combined with the known c_t values. The optimum value of y can thus be found without knowing the optimum values of the decision variables. If the minimum total cost, y^* , is acceptable, Equation 2.5 can be used to find the optimum values of the decision variables. Thus, the total cost of several designs can be compared before the values of the individual decision variables are obtained.

When applied under favorable conditions, geometric programming

exhibits remarkable ability to locate optima. For illustrative purposes, consider the following unconstrained optimization problem. In this simple example problem, it is desired to find the dimensions of a cylindrical storage tank which would minimize the construction and operating costs for the tank and filling system. Let the total cost be denoted by C_T

$$C_T = c_1 \frac{k}{\frac{\pi d^2 h}{4}} + c_2 \frac{\pi d^2}{4} + c_3 \pi d h \quad (2.11)$$

where d and h are the diameter and height of the tank and c_1 , c_2 , and c_3 are, respectively, the unit costs of filling the tank, material for the base, and material for the sides. The quantity k is the total volume to be supplied from storage in a period of time such as a year. In this crude example, the operating costs are represented by the first term while the construction cost is represented by the latter two terms.

Grouping the constants together, Equation 2.11 can be rewritten in the following form

$$\begin{aligned} C_T &= (c_1 \frac{4k}{\pi}) d^{-2} h^{-1} + (c_2 \frac{\pi}{4}) d^2 + (c_3 \pi) dh \\ &= c_1' d^{-2} h^{-1} + c_2' d^2 + c_3' dh \end{aligned} \quad (2.12)$$

Equation 2.12 is similar in form to Equation 2.1, and, as all $c_t' \geq 0$, geometric programming may be used to solve the problem. Applying Equations 2.6 and 2.7, one obtains

$$\begin{array}{rcccccl} w_1 & + & w_2 & + & w_3 & = & 1 \\ -2w_1 & + & 2w_2 & + & w_3 & = & 0 \\ -w_1 & & & + & w_3 & = & 0 \end{array}$$

Solving the equations above for the weights, the following results are obtained.

$$w_1 = 2/5, \quad w_2 = 1/5, \quad w_3 = 2/5$$

It is interesting to note that, without knowing the individual cost coefficients for filling the tank or constructing it, it is possible to learn that, at the optimum size, 40% of the total cost is incurred in filling the tank, 40% in constructing the sides, and 20% in constructing the base. This distribution of costs remains invariant regardless of the value of the c_t 's. These *invariant properties* are characteristic of many geometric programming solutions and provide additional insight to the analyst.

To evaluate the minimum cost, C_T^* , the cost coefficients and the weights are substituted into Equation 2.10. If the coefficients are assumed to be $c_1=10$, $c_2=6$, and $c_3=2$ and k is assumed to be 80×10^4 , the resulting C_T^* is found to be

$$C_T^* = \$5191$$

The values of the variables may be computed using any two weights of the type given in Equation 2.5. Using the term for w_2

$$(1.5)d^2 = (0.2)(5191)$$

$$d = 14.8$$

likewise using the third term

$$2dh = (0.4)(5191/\pi)$$

$$h = 22.3.$$

Constrained Optima

The general geometric programming problem is to minimize a posynomial subject to constraints of the form

$$y_m \leq 1, \quad m = 1, \dots, M \quad (2.13)$$

where y_m is also a posynomial. The constrained optima problem is solved by extending Equation 2.10 to include each term of the constraints. A weight for each constraint is added to the weights for each term as outlined elsewhere [19,56].

The ability to handle constraints makes geometric programming applicable to a wide range of engineering problems. Some constraints are natural forms of the problem, whereas others result from the process of converting functions to posynomial form. For example, the function

$$y(x) = (x_1^{-1} + x_2)^a (x_1 + x_2^{-2})^b$$

cannot be minimized directly since it is not a posynomial. However, it can be replaced by the equivalent problem

$$\text{Minimize } y(x) = x_3^a x_4^b$$

$$\text{Subject to } x_1^{-1} x_3^{-1} + x_2 x_3^{-1} \leq 1$$

$$x_1 x_4^{-1} + x_2^{-2} x_4^{-1} \leq 1$$

This new problem falls into the framework of geometric programming.

Degrees of Difficulty

Equation 2.6 and 2.7 provide $N+1$ linear dual variables, where N is the number of primal variables. There are, however, T dual variables in the dual equations, where T is the number of primal terms. The dual function is then unique if

$$T = N+1 \quad (2.14)$$

$$\text{If } T < N+1 \quad (2.15)$$

the primal problem is generally unconstrained and therefore not considered by geometric programming.

The degree of difficulty, D , is defined as

$$D = T - (N+1) \quad (2.16)$$

As D increases, the number of independent dual variables (equal to D) also increases. The dual problem is then a function of D independent and $N+1$ dependent variables.

If an optimum point exists, the minimum of the primal problem equals the maximum of the dual. The task of finding this point is a numerical problem increasing in difficulty as D gets larger, thus explaining the meaning of the term "degree of difficulty".

$$C_T = (c_1 4K) / (\pi d^2 h) + (c_2 \pi d^2) / 4 + c_3 \pi d h + c_4 h \quad (2.17)$$

with $c_4 = 1$. If, as in the previous example, this objective describes the cost of building and operating a water storage tank, the additional term could be considered the cost of adding supports to the sides. Using geometric programming, the weights are introduced following Equations 2.6 and 2.7.

$$\begin{aligned} w_1 + w_2 + w_3 + w_4 &= 1 \\ -2w_1 + 2w_2 + w_3 &= 0 \\ -w_1 + w_3 + w_4 &= 0 \end{aligned}$$

The minimum cost C_T^* cannot be computed until the optimum weights are known. Using the three equations above, it is possible to express each of the weights in terms of any one of the weights. Expressing all in terms of w_3 , the following are obtained.

$$w_1 = \frac{1}{3} (1 + \frac{1}{2} w_3); w_2 = \frac{1}{3} (1 - w_3); w_3 = w_3; w_4 = \frac{1}{3} (1 - \frac{5}{2} w_2)$$

Any choice of w_3 will produce values for the other weights that will satisfy Equations 2.6 and 2.7. The problem is now to select the optimum weights from among those possible using Equation 11.

Equation 10 holds for any number of degrees of difficulty, and the minimum cost will be given as a function of w_3 alone. This equation will be denoted as $d(w_3)$, the *substituted dual function*.

$$d(w_3) = \left(\frac{4Kc_1}{\pi(1/3(1+1/2w_3))} \right)^{\left(\frac{1+1/2w_3}{3} \right)} \left(\frac{\pi c_2}{4[1/3(1-w_3)]} \right)^{\left(\frac{1-w_3}{3} \right)}$$

$$\left(\frac{c_3 \pi}{w_3} \right)^{w_3} \left(\frac{c_4}{[1/3(1-5/2w_3)]} \right)^{\left(\frac{1-5/2w_3}{3} \right)}$$

Duffin [19] has proved that the substituted dual function is maximized by the optimum weight w_3^* and that this maximum value is equal to the minimum cost C_T^* .

$$C_T^* = d(w_3^*) = \max_{w_3} [d(w_3)]$$

In general, Duffin showed that a sufficient condition for C_T to be minimized is that the dual function defined by

$$d(w_1, \dots, w_T) = \prod_{t=1}^T \left(\frac{c_t}{w_t} \right)^{w_t}$$

be maximum with respect to w_t , subject to the conditions given in Equations 2.6 and 2.7.

The numerical value of w_3^* can be found by obtaining the first derivative of the logarithm of the substituted dual function and setting this derivative equal to 0. When this is done, the optimum value of w_3 is found to be

$$w_3^* = 0.395$$

The remaining optimum weights can be found using the relations shown in Equation 11.

$$w_1 = 0.399; \quad w_2 = 0.202; \quad w_3 = 0.395; \quad w_4 = 0.004$$

The weights can then be used as before to find the optimum values of d , h , and C_T^* .

$$d = 14.94 \quad h = 21.94 \quad C_T^* = 5213$$

Computer Algorithm

Edwards [20] has developed a computer program, which he calls the Geometric Programming System (GPS), to solve generalized geometric

programming problems. GPS with its eight subprograms is written in Fortran IV for the IBM 360/65 at Texas A&M University. If degrees of difficulty exist in the problem to be solved, the program calls a "pattern search" algorithm developed by Letto [37] to maximize the dual function. GPS provides the following output:

1. The primal problem.
2. The dual matrix.
3. The value of the dual function and dual variables.
4. The solution to the primal problem including the optimum value of primal variables and the value of each term and each row.
5. The absolute and relative difference between the primal and dual functions.
6. The elapsed time used to solve the problem.

The next chapter describes the computer program in more detail.

Generalized Polynomials

Passy and Wilde [47] extended the original posynomial form to permit negative coefficients and reversed inequalities (constants ≥ 1). This extension, however, releases the problem from its convexity form permitting the location of stationary points but no longer guaranteeing a global minimum.

Consider as the generalized polynomial

$$y_m(x) = \sum_{t=1}^{T_m} \sigma_{mt} c_{mt} \prod_{n=1}^N x_n^{a_{mnt}} \quad (2.18)$$

$$m = 1, \dots, M.$$

where $\sigma_{mt} = \pm 1$ to generalize the sign of c_{mt} ,
and all $c_{mt} > 0$. T_m is the number of terms in y_m .

The primal problem is then to

Minimize $y \equiv y_1^*$

subject to the constraints

$$y_m \leq \sigma_m (= \pm 1), m = 2, \dots, M.$$

and

$$x_n > 0, n = 1, \dots, N.$$

The dual problem contains T variables ω_{mt} satisfying a revised normality condition

$$\sum_{t=1}^{T_1} \sigma_{1t} \omega_{1t} = \sigma_1 (= \pm 1), \quad (2.19)$$

N revised orthogonality conditions:

$$\sum_{m=1}^M \sum_{t=1}^{T_m} \sigma_{mt} a_{mnt} \omega_{mt} = 0, \quad n = 1, \dots, N, \quad (2.20)$$

T nonnegativity conditions

$$\omega_{mt} \geq 0, \quad m = 1, \dots, M$$

$$t = 1, \dots, T_m$$

and $M - 1$ inequality constraints

$$\omega_{m0} \equiv \sigma_m \sum_{t=1}^{T_m} \sigma_{mt} \omega_{mt} \geq 0, \quad (2.21)$$

$$m = 2, \dots, M$$

where

$$T = \sum_{m=1}^M T_m,$$

and σ_1 can be shown to be of the same sign as y^* . The reader will note that the weight are designated ω_{mt} in the generalized case to differentiate it from the posynomial case.

Then the dual function is constructed as follows:

$$d(\omega) \equiv \sigma_1 \left[\prod_{m=1}^M \prod_{t=1}^{T_m} \left(\frac{c_{mt} \omega_{m0}}{\omega_{mt}} \right)^{\sigma_{mt} \omega_{mt}} \right]^{\sigma_1}, \quad (2.22)$$

where

$$\omega_{10} \equiv +1,$$

and, by L'Hôpital's rule,

$$\lim_{\omega_{mt} \rightarrow 0} \left(\frac{c_{mt} \omega_{m0}}{\omega_{mt}} \right)^{\sigma_{mt} \omega_{mt}} = 1.$$

If all $\sigma > 0$, the primal problem is convex in its e-transformation, and at the optimum

$$y_1(\underline{x}) \geq y_1(\underline{x}^*) = d(\underline{\omega}^*) \geq d(\underline{\omega}). \quad (2.23)$$

In a nonconvex case, the local minima of the primal function correspond to the dual critical points such that

$$y_1(\underline{x}^0) \equiv \min y_1(\underline{x}) = \max d(\underline{\omega}) \equiv d(\underline{\omega}^0). \quad (2.24)$$

These points may be local maxima, minima, or saddle points.

After the proper values for the dual function are found, the primal variables can be solved for by the following equations, which are linear in $\ln x$ and are analogous to Equation 2.5 for the convex case.

$$c_{1t} \prod_{n=1}^N x_n^{a_{1tn}} = \omega_{01} \sigma_1 y_1^0, \quad t = 1, \dots, T_1 \quad (2.25)$$

$$c_{mt} \prod_{n=1}^N x_n^{a_{mnt}} = \frac{\omega_{mt}}{\omega_{m0}}, \quad t = 1, \dots, T_m \quad (2.26)$$

$m = 2, \dots, M.$

It can be seen from Equation 2.25 that σ_1 has the same sign as y_1^0 . There are $T_M > N$ equations, so N can always be found, and the value of each primal variable can be determined.

COMPUTER ALGORITHM

As a part of this research, a computer algorithm for solving geometric programming problems was developed. This program is discussed in this chapter. The notation employed is the $x - w$ notation for the generalized problem [2]. Since there are no lower case characters on the line printer, capital letters in the text notation are prefaced by CAP- in the computer code. (See Appendix A). Thus, M represents m , and CAPM means M . However TSUBM(M) stands for T_m .

$$\begin{aligned} \omega_{mt} &= w_i & t &= 1, \dots, T_m \\ & & m &= 1, \dots, M, \\ \{l &= t + \sum_{i=0}^{m-1} T_i\} \end{aligned} \quad (3.1)$$
$$\omega_{m0} = w_{T_M} + m + 1, \quad m = 0, \dots, M.$$

Main Program

GPS, the main program, and its eight subprograms have been written in Fortran IV for the IBM 360/65 at Texas A&M University. See Appendix D for a description of subroutine interfaces.

GPS (for Geometric Programming System) reads in data and prints out the primal problem. (See Figure 3.1). The augmented matrix of coefficients of the dual variables is constructed in array B, and subroutine MAT is called to solve the system, if possible, or to place an identity matrix in the left hand side if it is not. From Equations 2.19 and 2.20, it will be seen that the system of dual equations occupies a matrix of $N + 1$ rows and T_M columns. Thus if degrees of difficulty exist, that is, $D \equiv T_M - (N + 1) > 0$, there will be D columns on the right side of the matrix which represent independent dual variables. The $T_M + 1$ st column is the constant vector in the augmented matrix which is printed.

If degrees of difficulty exist, a pattern search algorithm by A. R. Letto [37] is called to maximize the dual function as a function of the D independent dual variables, constrained by the natural nonnegativity requirement on the $T_M + M$ dual variables. Subroutine VAL is then called to evaluate the dual function at this maximum point, and the dual solution is printed.

From Equations 2.25 and 2.26, it has been shown that the equations involving the primal variables are linear in $\ln x$. Matrix B is loaded with the linearized equations, and once again MAT is called to solve the system. This time the matrix has T_M rows and N

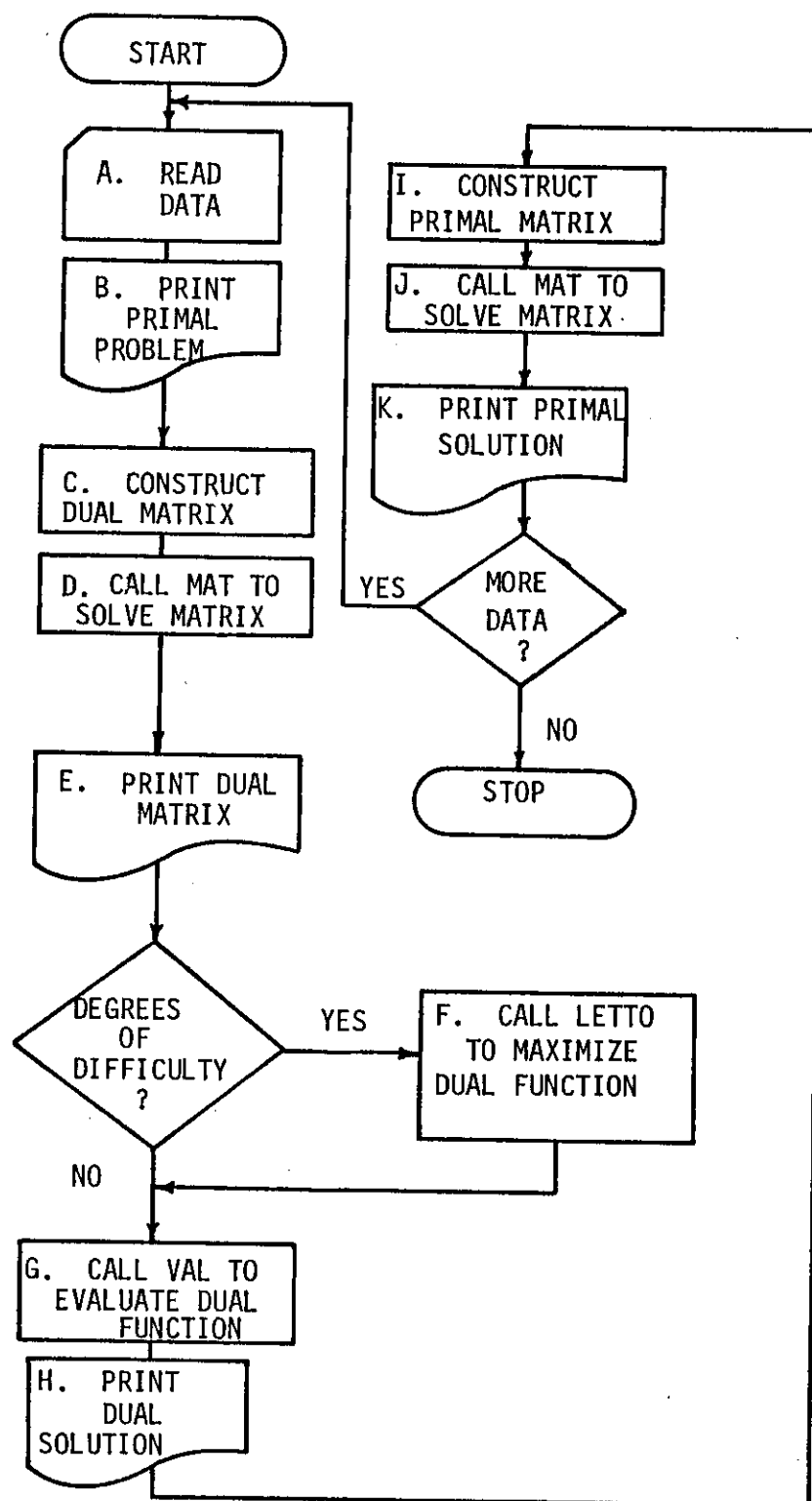


Figure 3.1 Program GPS.

columns, $T_M > N$. However, it is now known whether or not all the primal variables appear in the top N rows; so the routine solves for them wherever they appear. The primal solution is then printed, and, if another problem does not follow, execution is terminated.

Subroutine MAT

MAT is called twice, once to solve the dual matrix, then to solve the primal matrix. The general characteristics of the routine are noted above. The method of solution is simply algebraic addition of equations, after the method of Gauss-Jordan.

The pivot element (see Figure 3.2) referred to is simply the element on the upper left to lower right diagonal. The subtraction of rows is performed in the usual weighted manner, so that the elements in the pivot element's column are summed to zero.

The algorithm above (10) forms B into an upper-triangular matrix, and below (10) it finishes the job of creating the left side identity matrix. Note that the exchange of rows and columns, which always occurs downward and to the right, effectively insures that no information concerning the solution shall be lost if it is sought.

Subroutine LETTO

When the dual matrix has been constructed, if degrees of difficulty exist, LETTO is called. The principal elements given to LETTO are the dual matrix B and the number of degrees of difficulty. The D independent dual variables become the pattern search variables, and all the dual variables are invoked as constraints (see Output). LETTO calls VAL, SUMS, SUMC, and OUTPUT. It manipulates the D independent

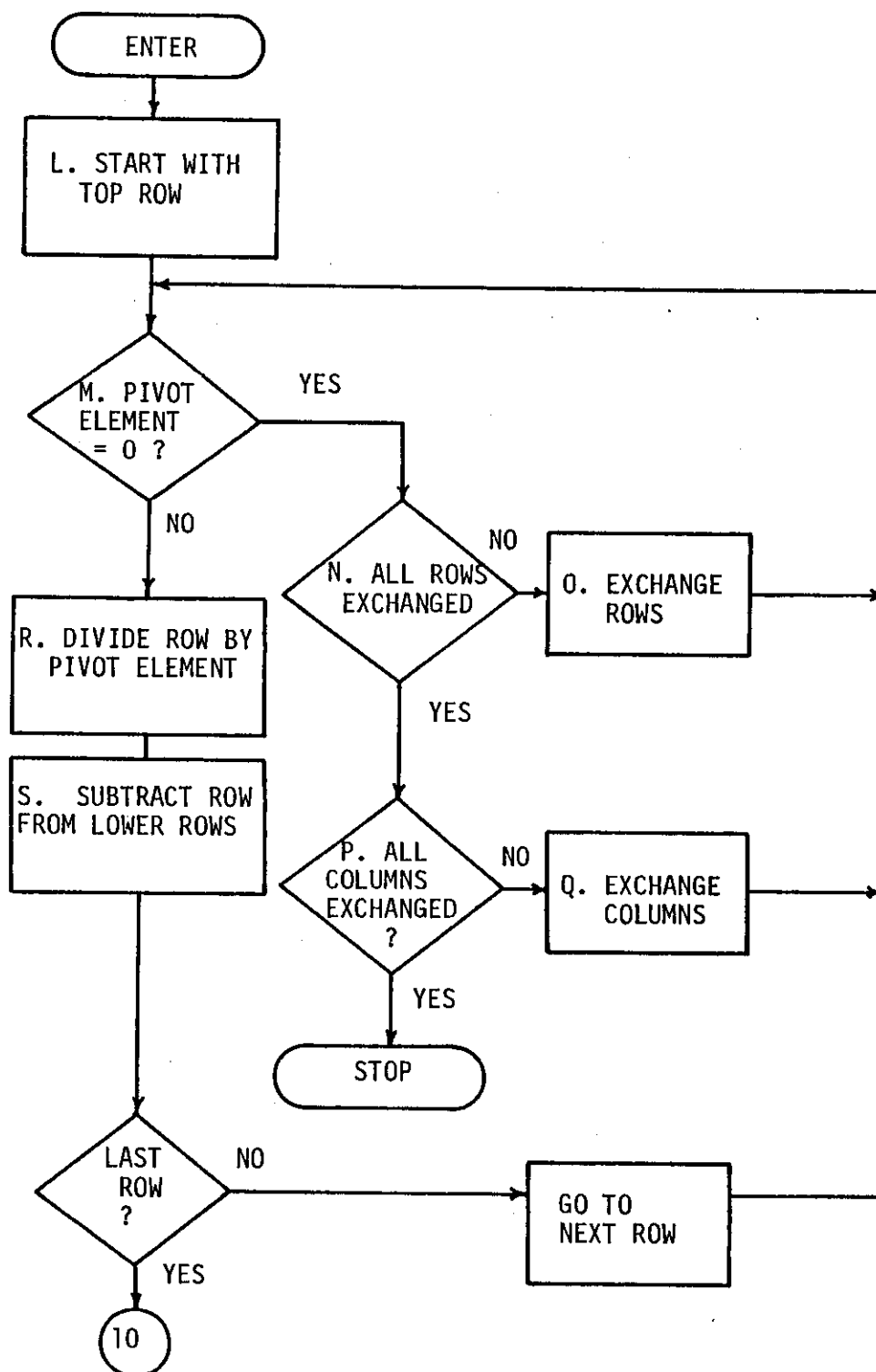


Figure 3.2 Subroutine MAT Flow Diagram

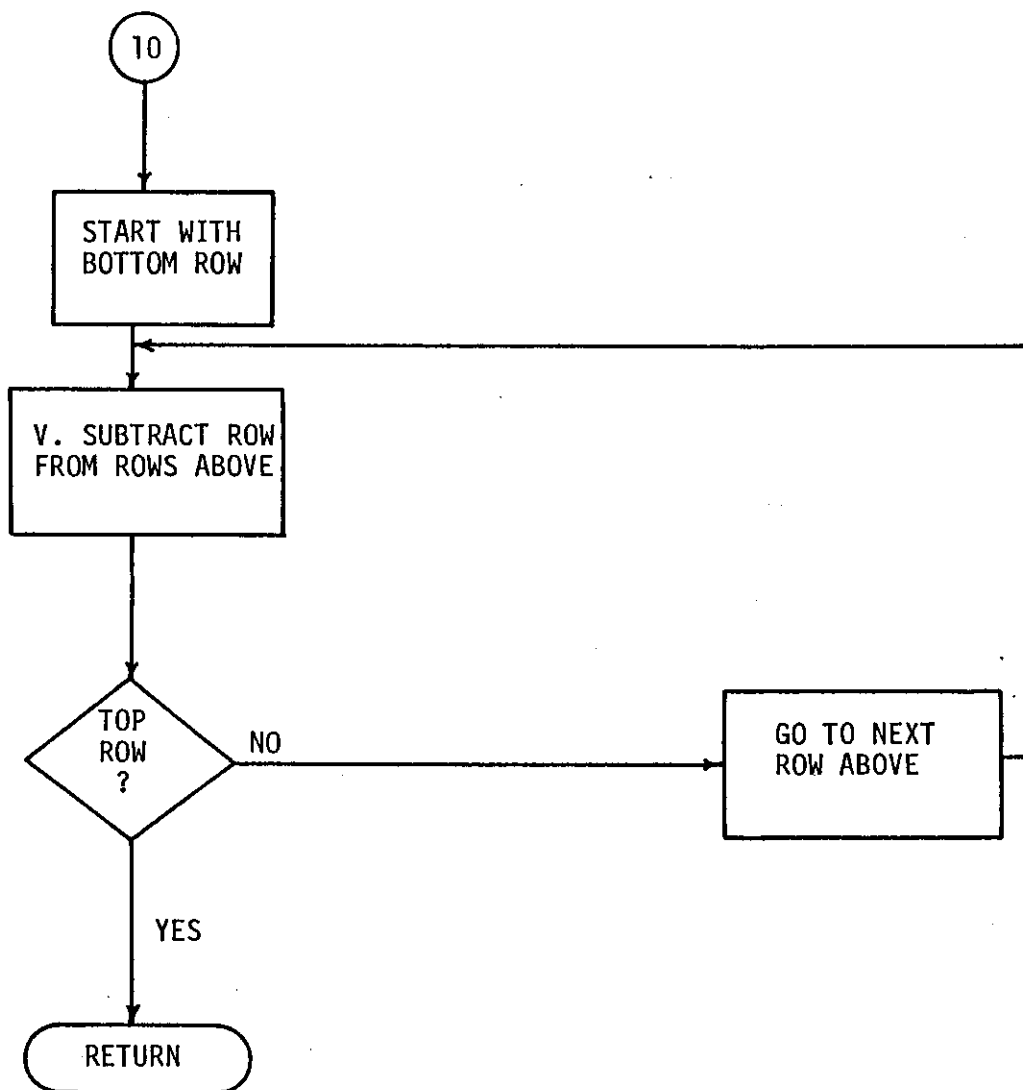


Figure 3.2. Continued.

dual variables and, evaluating the dual function and constraints through VAL, proceeds to search out a maximum of the dual.

The search method is after R. Hooke and T. A. Jeeves [25] and basically involves two kinds of moves. The first is exploratory search, in which each variable is perturbed slightly from its starting point. If no improvement is noted in the dual function, the variables are perturbed with a smaller step size. This may be repeated until the step size falls below some lower limit, when the search is successfully terminated. In the printout, this is called "local search." If the perturbation of any variable produces an improvement (larger value) over the base point, the remaining variables are perturbed around this new base point. The difference between the final and original base points defines a vector which lies "uphill" and approximates the route of steepest ascent (direction of largest gradient).

The second kind of move, the pattern move, proceeds in the direction of this vector. Each step goes twice the distance of the preceding step. This multiplication factor gives the method its power. When the last point found is worse than the previous point, the method reverts to exploratory search about the previous point. Thus the method never proceeds to a worse point.

When a constraint is contacted (some $\omega < 0$), a gradient summation technique is employed to enable the routine to proceed uphill along the edge of the constraint. If the starting point is infeasible, the sum of all the constraints becomes the objective

function to be maximized. Feasibility is reached when all constraints are positive. The Letto algorithm performs rapidly and seems always to find an optimum if one exists, and a global optimum if the region is concave.

Subroutine VAL

VAL receives the values of the independent dual variables, if there exist degrees of difficulty, and solves for all the dual variables from the matrix B, which has been simplified so that the calculation can be performed rapidly. This subroutine then calculates the value of the dual function, knowing \bar{g} , $\bar{\omega}$, and all c_{mt} .

Other Subroutines

Subroutine BLOWUP is called whenever an obvious error in data is read. It explains the error and terminates execution, since the next problem could not be read in while data for the present problem is still in line.

Subroutine OUTPUT prints out the values of the dual function, independent dual variables, and all constraints at each step in the pattern search. The constraints are in fact all the $\bar{\omega}$ in order. (See Equation 3.13). $\omega_{10} = w_{T_M+1} \equiv 1$ is included for consistency.

The BLOCK DATA subroutine initializes certain variables. Note that \bar{x} is initialized to 0.2 for the first problem, but is never reinitialized, except for $\bar{x}_1 = 0$ for any zero degree of difficulty problem. This is done so that similar problems can be read in following one another, and the time required to find the close optimum in the pattern search will be small. Indeed, except in such cases,

any value for λ , $0 < \lambda < 1$, is about as good as any other, since it would be very difficult to pre-determine the dual optima.

Input

The program will handle generalized polynomials for which
 $M \leq 10$; $N \leq 10$; $T_m \leq 10$, $m = 1, \dots, M$; $D \leq 20$.

The input consists of:

Card 1; (22I2 format)

$M, N, (T_m, m = 1, \dots, M), (\sigma_m, m = 1, \dots, M)$

Card 2; (A80 format)

Title of problem

Card 3; (E 10.0, I1 format)

XLLIM, the lower limit on variable moves in the Letto
 pattern search routine, and

LPRINT. If LPRINT = 0, the pattern search values are
 printed out at each step, with all messages.

If LPRINT = 1, these values are printed out only
 at beginning and end of the Letto routine. All
 messages are printed out at each step.

If LPRINT = 2, only feasibility messages are
 printed out.

If LPRINT = 3, only final values are printed out.

The next cards read in the coefficients and exponents of each
 term of the primal function and all constraints in order. For
 example, the first two terms of the primal function might be:

$$c_{11} x_1^{a_{111}} x_2^{a_{112}} x_3^{a_{113}} + c_{12} x_3^{a_{123}} x_5^{a_{125}} x_2^{a_{122}}$$

Then the cards would be as follows:

Card 4 (E 25.0 format) (If E is punched, it must be right justified.)

c_{11}

Cards 5-7 (I2, E 16.8, I2 format)

1, a_{111} , 0

2, a_{112} , 0

3, a_{113} , 1 (1 denotes last variable in term)

Card 8 (E 25.0 format)

c_{12}

Cards 9-11 (I2, E 16.8, I2 format)

3, a_{123} , 0

5, a_{125} , 0

2, a_{122} , 1

and so forth to the last term in the last row of the generalized polynomial.

If there is another problem to be run at this time, succeeding

cards should follow the same order as before, starting with Card 1.

After the last problem, one final card should contain 1 in column 2 and otherwise be blank.

Output

A typical output is given in Appendix B. The line below the main heading is the title of the problem, read in on the second card. The generalized polynomial is then printed out, with the horizontal lines separating the objective function and constraints. The symbol .L.E. on the right implies that the constraint $\leq \sigma$ ($\equiv \pm 1$).

On a new page the augmented dual matrix is printed out after having the identity matrix placed in its left side. The rightmost column is the constant or result column. The next D columns to the left represent the independent dual variables.

The pattern search output is begun if $D > 0$. The print option determines how much is printed out.

The value of the dual function and of all the dual variables is printed out next. Then the primal solution, including the optimal values of the variables and the values of each term and each row of the generalized polynomial primal problem, evaluated at those values, is printed out.

The absolute and relative difference between the primal and dual functions is shown. If the problem is convex, the optimal value can be shown to lie between the two functional values; *i.e.*, an upper and lower limit and maximum error is determined. The elapsed time finishes the print-out.

Error Messages

The error messages from subroutine BLOWUP are printed out after the data in question is printed, and should be self-explanatory.

One message from MAT can be printed to signify that the matrix in question is singular. That is, there are zeros in all of the elements to the right and below a certain pivot element, and thus the matrix cannot be solved. If the message occurs after the primal problem is printed out, the singularity was found in the dual matrix. If it occurs later, after the dual variables are printed out, the fault lies in the primal matrix. Job execution stops, but the next problem may be read in.

The message $SIG(1) = -SIG(1)$ signifies that the program has realized that the sign guessed for the objective function at a minimum was incorrect. The value is changed and execution continues without prejudice to the results.

CHAPTER IV

WATER FILTRATION THEORY

The Filtering Process

The purpose of this chapter is to provide a basis for the model development associated with the comprehensive example optimization problem involving water filtration which is described in Chapter V. The solution of this problem is discussed in Chapters V and VI.

Water filtration is the process of passing water through a porous media to remove turbidity, suspended matter, and some colloidal material. The turbidity and suspended matter may be present in the raw water or may result from some precipitation action preceding the filtration process in the water treatment sequence. Particle removal in filtration involves two mechanisms -- a physical transport mechanism and a chemical attachment mechanism. Most filtration theory and all of the mathematical models are based on the physical mechanism [4,11,31,32,29]. The chemical mechanism has been suggested [44], but it has not been mathematically modeled. A review of some of the important literature follows.

Physical Removal Mechanisms

Iwasaki, [33] the founder of modern filtration theory, developed three fundamental equations describing the filtration process

$$-\frac{\partial C}{\partial L} = \lambda C \quad (4.1)$$

$$\lambda = \lambda_0 + c\sigma \quad (4.2)$$

$$\frac{\partial C}{\partial L} + \frac{\partial \sigma}{\partial L} = 0 \quad (4.3)$$

in which, C = concentration of suspended particles, L = filter depth, λ = coefficient of proportionality (also termed the filter coefficient), σ = volume of suspended matter retained per unit volume of filter, and c = a constant.

Equation 4.1 states that the amount of matter arrested in a lamina is proportional to the amount that enters the lamina. Equation 4.2 states that the change in the proportionality factor or filter coefficient is directly proportional to the density of the matter filling the voids. Equation 4.3 is a mass balance or continuity equation stating that the decrease of flow through a lamina is equal to the increase of the density therein.

Ives [31] modified Equation 4.2 stating that the filter coefficient, λ , first increases linearly with the specific deposit, σ , and then decreases nonlinearly.

$$\lambda = \lambda_0 + k\sigma - \frac{\phi\sigma^2}{p_0 - \sigma} \quad (4.4)$$

Ives also modified Equation 4.3 to

$$-\frac{\partial C}{\partial L} = \frac{1-f_\sigma}{v} \frac{\partial \sigma}{\partial t} \quad (4.5)$$

and added a new equation to describe the headloss.

$$\frac{\partial h}{\partial L} = \left(\frac{dh}{dL} \right)_0 + k\sigma \quad (4.6)$$

where λ_0 = initial filter coefficient, p_0 = porosity of clean bed, ϕ = a filter parameter, t = filtration time, f_0 = self porosity of deposited solids, h = headloss, k = headloss constant, v = approach velocity, and $(dh/dL)_0$ = initial (clean bed) headloss per unit depth.

Others including Sholji [51] and Deb [11] have suggested modifications to the work of Iwasaki and Ives, but the basic principles have not been changed. Deb [10] used a computerized finite-difference method to obtain a numerical solution to his resulting nonlinear partial differential equations. Deb's model is, however, restricted to homogeneous, unisize, granular media which somewhat limits its practical value.

Camp [4] took another approach to developing a mathematical model combining a modified form of the Kozeny equation of laminar flow [34] with Stein's refinement [52] of the Iwasaki model. Camp suggested two hydraulic gradient equations.

$$i = \frac{k(1-p_0+\sigma)^2}{(p_0-\sigma)^3(d+\Delta d)^2} \quad (4.7)$$

$$i = \frac{k(1-p_0)^2}{d_0^2 p_0^2 (p_0-\sigma)} \quad (4.8)$$

where i = hydraulic gradient (dh/dL), d_0 = grain size at the start of filtration, Δd = increase in grain size at a particular time during a filter run in which the value of k is known, $k = (j u/g)v$, j being an overall characteristic of the filter bed media, u is the

kinematic viscosity of the fluid, g is the gravitational constant, and v is the approach velocity.

Equation 4.7 applies where the clogging or removal mechanism is due to the formation of a sheath on each grain. Equation 4.8 describes the situation where the clogging phenomena is due primarily to straining.

Laboratory filtration experiments [45,30] designed to compare the various equations of Ives, Camp, and others have shown a wide variance in effluent quality predictions. Such differences are probably due in part to the failure of the models to include the chemical mechanisms.

Chemical Removal Mechanisms

O'Melia and Struim [44] have written a comprehensive review of chemical filtration mechanisms. These mechanisms are usually related to Van der Waals forces and electrokinetic effects, but they have not been developed to a state where they can be expressed mathematically.

Van der Waals forces are the most powerful of the known attraction forces at the molecular level. These forces vary inversely at the seventh power of their difference of separation [9]. Electrokinetic effects are the result of an electric charge on the surface of colloidal particles caused by ionization or adsorption. Ions are centered around the solid particle in alternately charged layers. Electrokinetic repulsion between the outer layers tends to keep the particles dispersed. The presence of a charge on the surface of the filter media may, therefore, largely determine the stability of the colloidal dispersion.

In an attempt to include these chemical mechanisms with the physical, some researchers have posed a combination of theory and empirical results to describe the filtering process.

Empirical Studies

One of the more recent and comprehensive combinations of theory and empirical results was prepared initially by Hsiung for single media filters [26,27] and later extended to multimedia filters [8]. Hsiung proposed a new filter performance prediction theory based on a "random walk" analogy. If P represents the probability of penetration by a particle, the probability of a particle moving down to depth L at time t will be

$$P(L:rt,p) = \binom{rt}{L} p^L q^{(rt-L)} \quad (4.9)$$

where r = average rate of particle movement, p = probability of a success (particle moves one unit step downward), and q = probability of failure (the particle does not move).

The total number of particles arriving at L is

$$\partial m \sim P(L:rt,p) Q C_0 \partial t \quad (4.10)$$

Then by Iwasaki's continuity equation, Equation 4.3, modified to let $\partial L = Q \partial t$

$$\frac{\partial C}{\partial L} Q \partial t \sim P(L:rt,p) Q C_0 \partial t \quad (4.11)$$

or

$$\frac{1}{C_0} \frac{\partial C}{\partial L} \sim P(L:rt,p) \quad (4.12)$$

Thus particles are removed per unit depth of filter according to some probability function. Hsiung used a chi-square distribution to map this probability and proposed a deposit index, U , to translate filter performance data into a practical design parameter.

Hsiung then combined his theoretical aspects of filter performance with experimental data obtained from pilot plant studies for various grain sizes, flow rates, and influent concentrations. In the pilot plant studies, he first evaluated the effect of each variable on the filter performance for a given system holding the other variables constant. He then expressed this relationship in two performance curves [26,27]. The first curve compares U/L with $Q^a d^b t/L^c$, and the second compares $d^i (H_t - H_0)/Q^j C_0^k L^n$ with $Q^a d^b t/L^c$ in which, L = depth of bed, Q = filtration rate per unit area of filter, d = mean grain size of filter media, t = filtration time, H_t = headloss at time t , H_0 = headloss at beginning of filtration, C_0 = concentration of particles in influent, and a, b, c, i, j, k , and n are the parameters determined from the pilot plant studies. Plotted on log-log paper, both curves are concave. These performance curves can be developed for single or multi media filters. For multimedia, the equivalent grain size, d_e , is approximated by

$$d_e = x_1 d_1 + x_2 d_2 + x_3 d_3 \quad (4.13)$$

in which, d_1 , d_2 , and d_3 are the mean grain sizes of each media (assuming here that there are three sizes), and x_1 , x_2 , and x_3 are the percentage by volume of each media. A representative set of performance curves is shown in Figures 4.1 and 4.2. These curves were developed by Conley and Hsiung [8] to show a multimedia filter's typical pattern of performance for a synthetic muddy water after coagulation and settling. A secondary flocculant (polyacrylamide) was also applied immediately ahead of the filter. The addition of such polymers to the filter influent greatly increased filtration efficiency, but it also increased the head loss with resulting new performance curves. Thus, a filtration cost model developed around the performance curves cannot include the polymer as a variable since new constraints (based on the curves) would have to be used for varying amounts of polymer.

Multimedia Filters

Initially, filters were of the slow sand type with a filtration rate of 0.04 to 0.12 gallons per minute per square foot (gpm/sq ft). In the late 19th century, rapid sand filters were introduced operating at a rate of approximately 2 gpm/sq ft. Neptune-Microfloc has recently introduced a multimedia filter of anthracite, sand, and garnet in layers of decreasing size. Slow and rapid sand filters have their finer grains at the top and perform most of the filtration in the top few inches. The multimedia filter has its coarser grains at the top permitting filtration action through the entire depth of the filter (See Figure 4.3).

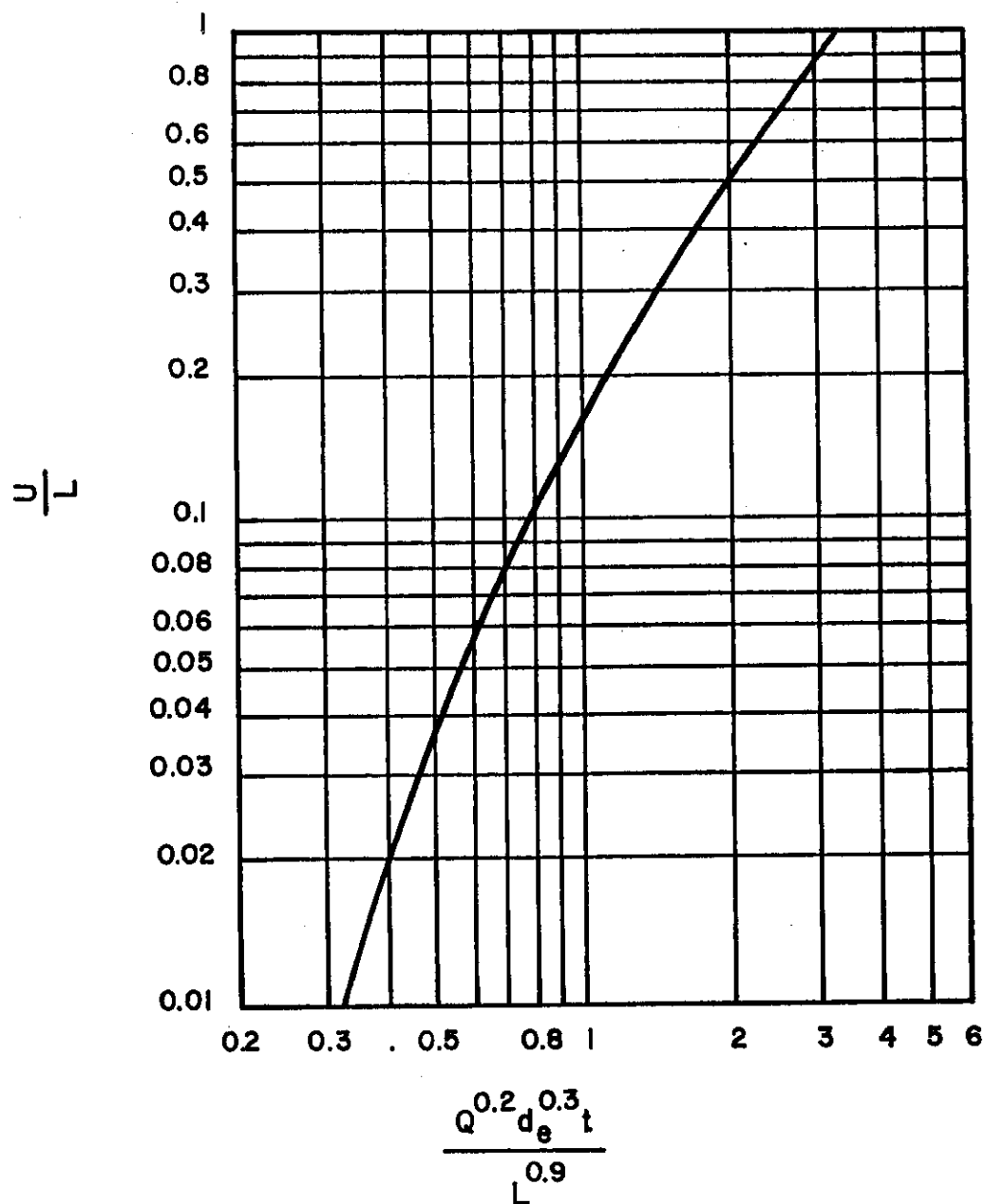


Figure 4.1 Performance Curve I for a Synthetic Muddy Water

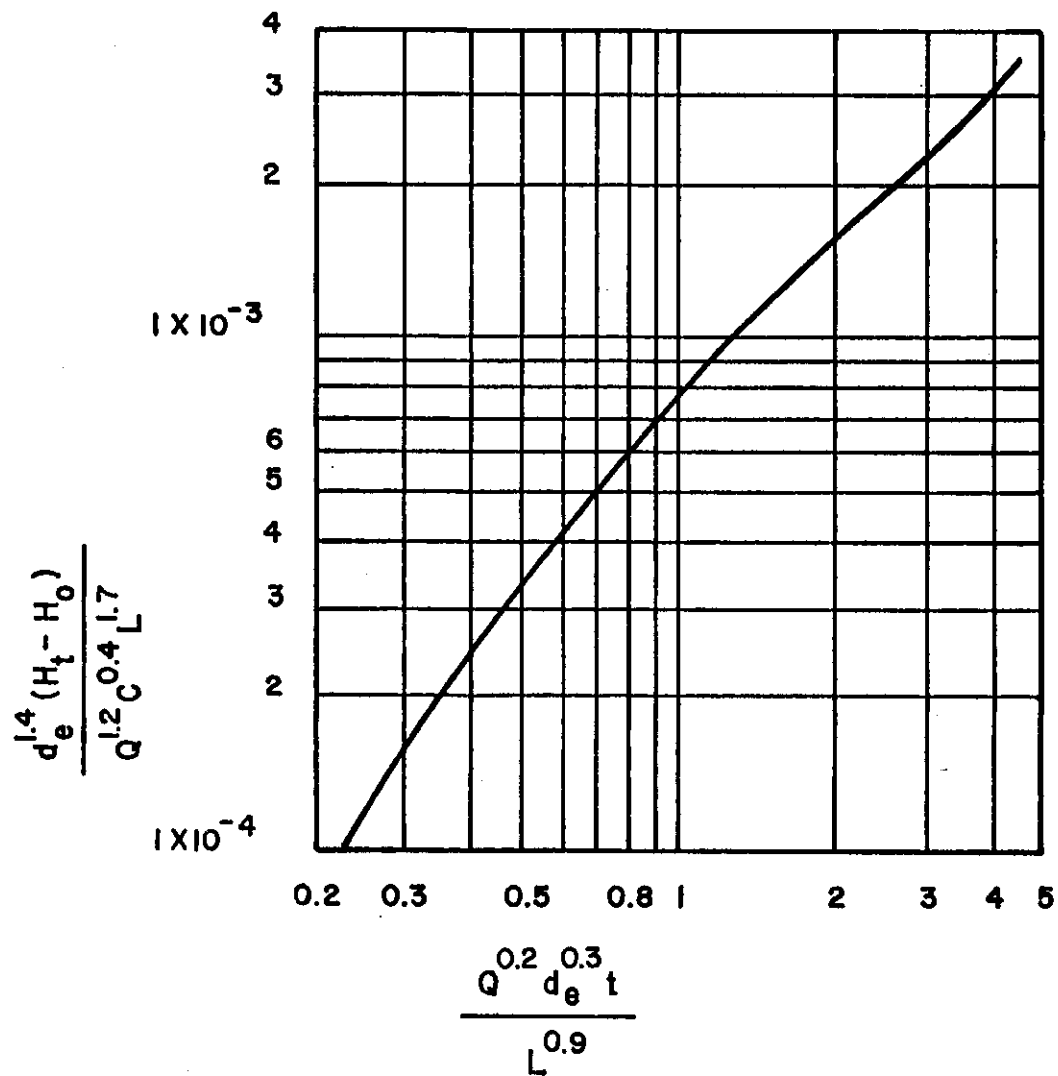


Figure 4.2 Performance Curve 2 for a Synthetic Muddy Water

1 ft.

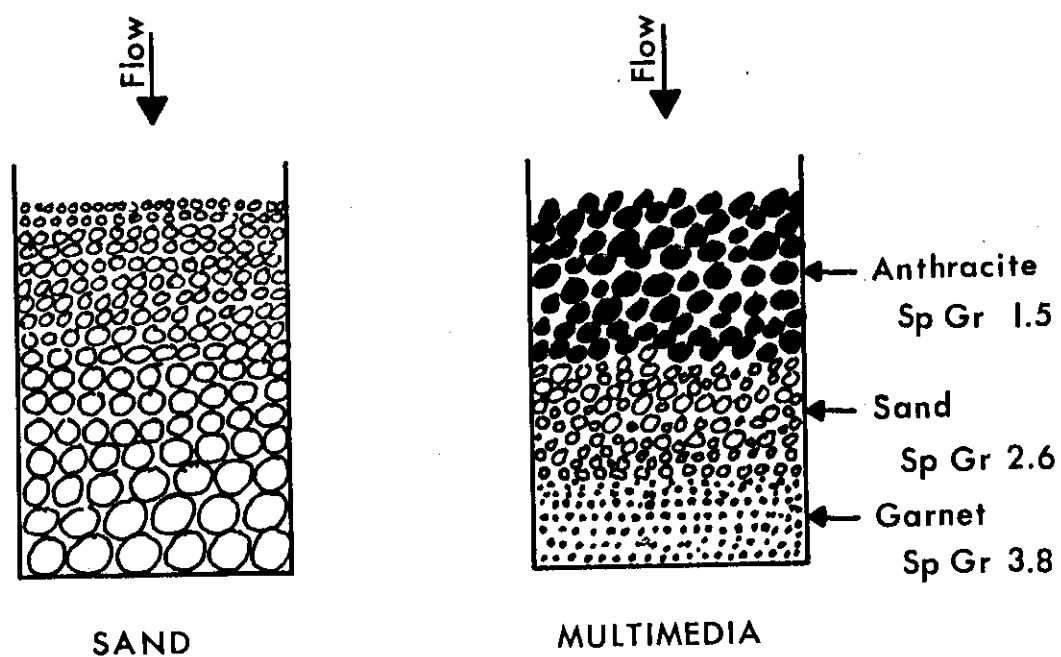


Figure 4.3 Comparison of Filter Media

Multimedia filters can operate at filtration rates up to 8 gpm/sq ft. These filters have gained wide acceptance [6,13,48,42] and offer an economical solution for treatment plants needing greater capacity than that of existing rapid sand filters [36,50]. In designing a multimedia filter, a wide range of choices are available for the specific gravity, size, and depth of each media layer. Selection of an optimum combination of media is dependent on several factors including turbidity loading, filtration rate, desired quality of effluent, and initial cost. In Chapter V, a cost model is developed to assist in the selection of an optimal media combination. First, however, the next chapter will review geometric programming, an optimization technique which will be used in finding a solution to the cost model.

CHAPTER V

FILTER COST MODEL

Reducing the Problem

There are a number of water filtration design problems to which a cost minimization technique could be applied. These problems range in scope from the design of a complete new facility to minor operating changes in an existing plant. A representative problem has been selected between these two extremes which adequately demonstrates the modeling technique and yet keeps the problem small enough to be tractable.

The problem chosen for presentation herein is the selection of an optimum trimedia to replace the rapid sand media in an existing filtration plant. A typical multimedia filter is shown in Figure 5.1. Such conversions have proven to be a very economical means of increasing plant capacity at minimal cost. For example, the water plant at Winnetka, Illinois, was converted from 6.0 to 12.0 million gallons per day (mgd) at less than 25 per cent of the cost of adding either new rapid sand or diatomaceous-earth filters [50]. Greenville, Texas, used multimedia filters to increase their plant capacity from 6.0 to 10.5 mgd at a cost of \$200,000 as compared to \$475,000 for new rapid sand filters [36].

In designing a multimedia filter, a wide range of choices are available for the size and depth of each media layer. Selection of an

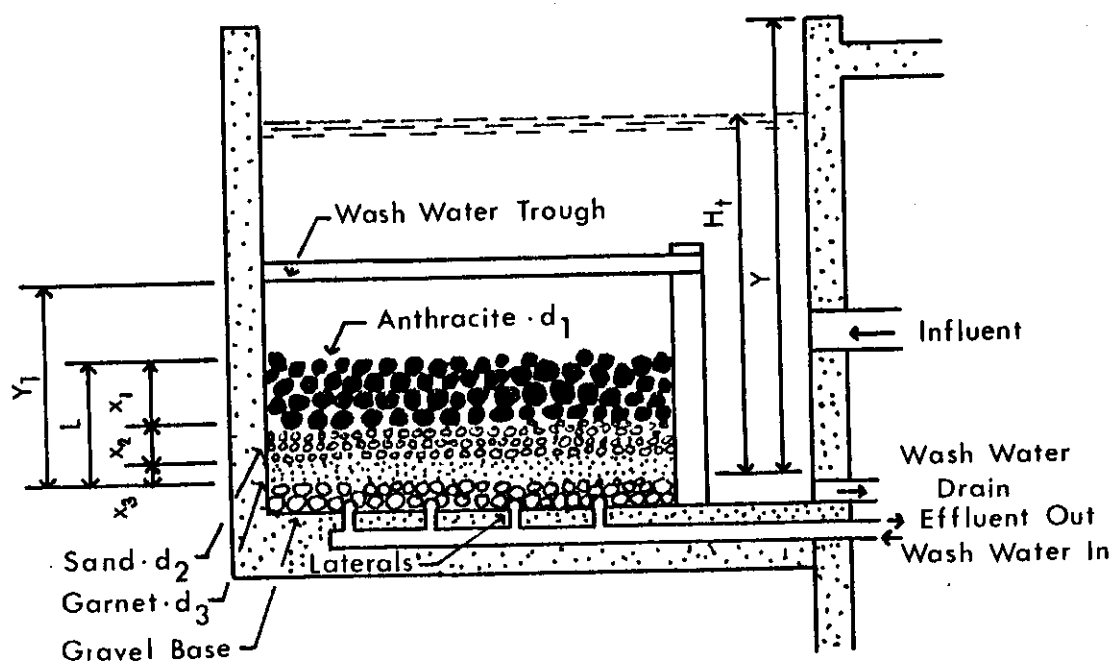


Figure 5.1 Multimedia Filter

optimum combination of media is dependent on many factors. These factors can be expressed in mathematical equations to form the objective function and constraints of the filter cost model.

Objective Function

The principal decision variables in the mathematical model are the depth of each media, flow rate, length of filter run, headloss, and the equivalent grain diameter. The diameter of each media could also be established as a decision variable with upper and lower bounds, but this greatly increases the number of variables requiring a large optimization program. To retain generality, the model will be presented in this chapter in its general form with all known variables including media diameter. Then in Chapter VI, some variables will be established as constants so that the Geometric Programming System (GPS) can be used to solve a sample problem.

The objective function in the mathematical model is designed to measure the effectiveness (in this case the cost) of media selection. The primary costs are the original cost of the media amortized over a specified period and the cost of the backwash water. Personnel operating costs, etc. are generally constant for all media and are therefore not included in the objective function. The objective function is expressed in dollars per million gallons based on a 30-day month average.

A series of payments which will repay the initial investment in

n years together with an annual interest of rate i on the uncovered portion can be found by [12]

$$\frac{i(1+i)^n}{(1+i)^n - 1} \quad (5.1)$$

Converting Equation 5.1 to a monthly series of payments and defining the resultant as R

$$R \equiv \frac{\frac{i}{12} (1 + i/12)^{12n}}{(1 + i/12)^{12n} - 1} \quad (5.2)$$

The objective function is then

$$\frac{\left[\frac{A R}{12} (x_1 P_1 + x_2 P_2 + x_3 P_3) + \frac{43,200 W A}{15t} \right] 10^6}{43,200 Q A} \quad (5.3)$$

in which, A = surface area of filter (sq ft); x_1, x_2 , and x_3 = depth of each media (in); P_1, P_2 , and P_3 = price of each media (\$/cu ft); t = time units of filtration period (15 min); W = backwash cost (\$/sq ft); the 43,200 represents the number of minutes in a 30-day month; and the 10^6 converts to million gallons.

Eliminating the A in the numerator and denominator and changing to posynomial form, Equation 5.3 becomes

$$(1.92)(R x_1 P_1 Q^{-1} + R x_2 P_2 Q^{-1} + R x_3 P_3 Q^{-1}) + 66,700 W t^{-1} Q^{-1} \quad (5.4)$$

The backwash cost, W , is equal to the product of the backwash rate, the backwash time, and the cost of backwash water.

There are several restrictions on the objective function's decision variables, Q , t , x_1 , x_2 , and x_3 . These restrictions can also be expressed in mathematical equations forming the constraints of the model.

Constraints

Before the constraints can be established, equations for the filter performance curves must first be developed. The two performance curves were previously described in Chapter IV and shown there as Figures 4.1 and 4.2. The concave curves can each be approximated by two straight line segments. An abscissa value of 0.9 has been selected for the break point for the straight line segments of each curve thereby establishing a uniform range of values for each set of performance equations.

The bottom portion of the curve in Figure 4.1 (abscissa $\leq .9$) has a slope of 2.36 while the top has a slope of 1.70. The line intercept is 0.01 at the ordinate and 0.03 at the abscissa. The equation of the bottom portion of the curve is therefore

$$\text{Log } \frac{U}{L} = \text{Log } (0.1) + 2.36 \text{ Log } \left[\frac{(Q)^2 (d_e)^3 (t)}{(L)^9} \right] - \text{Log } 0.3 \quad (5.5)$$

Taking the antilog

$$\begin{aligned} \frac{U}{L} &= .01 \left[\frac{(Q)^{.2} (d_e)^{.3} t}{.3(L)^{.9}} \right]^{2.36} \\ &= .172 (Q)^{.47} (d_e)^{.71} (t)^{2.36} (L)^{-2.12} \end{aligned} \quad (5.6)$$

In the geometric inequality form

$$5.82 (U) (Q)^{-.47} (d_e)^{-.71} (t)^{-2.36} (L)^{1.12} \leq 1 \quad (5.7)$$

Similarly, for the top portion

$$6.49 (U) (Q)^{-.34} (d_e)^{-.51} (t)^{-1.70} (L)^{.53} \leq 1 \quad (5.8)$$

For Figure 4.2, the slope of the bottom portion of the curve is 1.39, and the top slope is 1.0. The ordinate intercept is 10^{-4} and the abscissa intercept is 0.22. Thus the equation of the bottom portion of the curve is

$$\begin{aligned} \text{Log} \left[\frac{(d_e)^{1.4} (H_t - H_o)}{(Q)^{1.2} (c_o)^{.4} (L)^{1.7}} \right] &= \text{LOG} (10^{-4}) \\ + 1.39 \text{ Log} \left[\frac{(Q)^{.2} (d_e)^{.3} (t)}{(L)^{.9}} \right] &= \text{Log} (.22) \end{aligned} \quad (5.9)$$

Taking the antilog

$$\frac{d_e^{1.4}(H_t - H_o)}{(Q)^{1.2}(C_o)^{.4}(L)^{1.7}} = 10^{-4} \frac{(Q)^{.2}(d_e)^{.3}(t)^{1.39}}{.22(L)^{.9}} = 8.2 \times 10^{-4}(Q)^{.28}$$

$$(d_e)^{.42}(t)^{1.39}(L)^{-.125} \quad (5.10)$$

or

$$t^{1.39} = \frac{(H_t - H_o)(d_e)^{.98}}{8.2(10^{-4})(Q)^{1.48}(L)^{.45}(C_o)^{.4}} \quad (5.11)$$

Similarly the equation for the top portion of the curve in Figure 4.2 is

$$t = \frac{(H_t - H_o)(d_e)^{1.1}}{7.5(10^{-4})(Q)^{1.4}(L)^{.8}(C_o)^{.4}} \quad (5.12)$$

Equation 5.11 can be expressed in the geometric programming inequality form as

$$8.2(10^{-4})(Q)^{1.48}(L)^{.45}(C_o)^{.4}(d_e)^{-.98}(H)^{-1}(t)^{1.39} \leq 1 \quad (5.13)$$

and Equation 5.12 can be expressed as

$$7.5(10^{-4})(Q)^{1.4}(L)^{.8}(C_o)^{.4}(d_e)^{-1.1}(H)^{-1}(t) \leq 1 \quad (5.14)$$

where $H = H_t - H_0$ or $H + H_0 = H_t$. Thus

$$H H_t^{-1} + H_0 H_t^{-1} \leq 1 \quad (5.15)$$

For a minimum length of filter run, t_{\min}

$$t_{\min} t^{-1} \leq 1 \quad (5.16)$$

Equations 5.8 and 5.14 will be used initially for the filter performance constraints. After the model is optimized, the abscissa value of $(Q) \cdot^2 (d_e) \cdot^3 (t/L) \cdot^9$ will be calculated, and, if this value is less than or equal to 0.9, Equations 5.7 and 5.13 will be substituted for Equations 5.8 and 5.14 and the model reoptimized. Equations 5.8, 5.14, 5.15, and 5.16 are thus the first three constraints of the model. The variable d_e is calculated using Equation 5.13 as $d_e = X_1 d_1 + X_2 d_2 + X_3 d_3$ where d_i is the diameter size of media i and X_i is the percentage of the media by volume. The bed depth L is also a sum of the individual media layers, $x_i = X_i/L$.

$$x_1 L^{-1} + x_2 L^{-1} + x_3 L^{-1} = 1 \quad (5.17)$$

$$d_1 x_1 d_e^{-1} L^{-1} + d_2 x_2 d_e^{-1} L^{-1} + d_3 x_3 d_e^{-1} L^{-1} = 1 \quad (5.18)$$

Equations 5.17 and 5.18 must equal unity. If they are to be used in a geometric programming algorithm of inequalities, each will require two equations with one ≤ 1 and the other ≥ 1 .

Since the filter is to be installed in an existing bed, headloss must not exceed the available clearance, Y . The clearance is measured from the top of the filter walls down to the surface of the coarse gravel which is used in most multimedia filters to support the third media layer of fine garnet. That is

$$12H_t Y^{-1} \leq 1 \quad (5.19)$$

Another clearance restriction involves the wash water troughs. The expanded bed depth during backwash must not exceed this clearance or media will be lost with the wash water. The amount of expansion of each media during backwashing depends on the specific gravity of the media and the backwash rate. During the pilot plant studies (discussed in Chapter IV), an optimum backwash rate can be established and the percentage of expansion E_i determined for each type of media. The total expanded depth must not exceed the available clearance, Y_1 , of the wash-water troughs

$$x_1 Y_1^{-1} (1 + E_1) + x_2 Y_1^{-1} (1 + E_2) + x_3 Y_1^{-1} (1 + E_3) \leq 1 \quad (5.20)$$

To ensure that production exceeds expected demand, a minimum filter flow, G , in gallons per day (gpd) is established. This minimum production in turn sets a minimum on the filter flow rate with an allowance made for lost production time during backwash. It is estimated that at most four percent of the 1440 minutes in a day will be used for backwash. This percentage is based on three backwashes

per day with 15 to 20 minutes for each wash.

$$(.96)1440 Q A \geq G \quad (5.21)$$

or

$$.0007236 Q^{-1} A^{-1} G \leq 1 \quad (5.22)$$

A minimum bed depth is required to prevent unfiltered water from breaking through the bed. Hudson [28] concluded that breakthrough will occur when

$$\frac{Q d_{H_t}^3}{L} = B \frac{60}{T + 10} \quad (5.23)$$

in which, T = water temperature ($^{\circ}\text{F}$) and B = a breakthrough index which varies for different waters depending on the response to coagulation and degree of pretreatment. Assuming that d_e is the equivalent diameter for the trimedia

$$\frac{T + 10}{60} Q d_{eH_t}^3 L^{-1} B^{-1} \leq 1 \quad (5.24)$$

An upper bound d_{\max} and a lower bound d_{\min} on each media diameter can be established from information furnished by the media supplier. Thus

$$d_i d_{i_{\max}}^{-1} \leq 1, \quad i = 1, 2, 3 \quad (5.25)$$

and

$$d_i^{-1} d_{i \min} \leq 1 \quad i = 1, 2, 3 \quad (5.26)$$

For proper stratification of the media after backwash, a size ratio of adjacent media must be established. This ratio can be determined satisfactorily by the principle of equal settling [39]. The ratio has been calculated by Conley and Hsiung [8] as

$$\frac{d_i}{d_{i+1}} \leq \frac{(D_{i+1} - D_{mi})}{(D_i - D_m)} \frac{1}{1.6}, \quad i = 1, 2 \quad (5.27)$$

or

$$\frac{1.6 d_i (D_i - D_{mi})}{d_{i+1} (D_{i+1} - D_{mi})} \leq 1, \quad i = 1, 2 \quad (5.28)$$

where

$$D_{mi} = e_i D + 1/2(D_i + D_{i+1})(1 - e_i), \quad i = 1, 2 \quad (5.29)$$

in which, e_i = average porosity of the mixing interface (two adjacent media) of the expanded bed near the end of backwash, D_i and D_{i+1} = density of media to be compared, and D_{mi} = density of mixture.

Other constraints could possibly be developed on an individual basis for a specific design problem, but those presented in this

section are generally most restrictive on the decision variables.

An Optimum Solution

Equation 5.4 represents the objective function and Equations 5.8, 5.14, 5.15, 5.16, 5.17, 5.18, 5.19, 5.20, 5.22, 5.24, 5.25, 5.26, and 5.28 describe the constraints of the mathematical model.

A solution to the model can be found by completing the following steps:

1. Find values for the constants in the model.
2. Determine coefficients for each term in the objective function and the constraints and set up the mathematical form of the problem.
3. Optimize the model using GPS or some alternate procedure.
4. Using the results of the decision variables determined by step 2, calculate $Q \cdot d_e^2 \cdot t^3 / L^9$. If this value is less than or equal to 0.9, use Equations 5.7 and 5.13 for the first two constraints and return to step 2 above.

A sample set of constants will be identified in the following chapter, and an optimization will be performed to test the model.

CHAPTER VI

A SAMPLE PROBLEM

Constant Values

The first step in the problem solution is to identify the constant values. These values can be obtained from plant records, pilot studies, and design criteria for the filter bed where the trimedia is to be installed. The values listed below are representative of those for a typical plant. The source of the value is shown where applicable:

i = interest rate = 5 percent

n = number of periods = 20 years

A = surface area of filter = 242 sq. ft.

P_1 = price of anthracite [46] = \$0.60/cu. ft.

P_2 = price of sand = \$0.20/cu. ft.

P_3 = price of garnet [1] = \$3.00/cu. ft.

Backwash rate [7,43] = 15 gpm/sq. ft.

Length of backwash [7] = 6.7 min.

Price of backwash water = \$0.35/thousand gallons

C_o = influent turbidity = 20 Jackson Turbidity Units (JTU)

C = effluent turbidity = 0.2 JTU

H_o = initial headloss = 1.333 ft.

t_{min} = minimum length of filter run = 40 periods
(each period is 15 minutes or 10 hours total)

d_1 = diameter of anthracite [7] = 1.0-2.0 mm

- d_2 = diameter of sand [7] = 0.4-0.8 mm
- d_3 = diameter of garnet [7] = 0.2-0.4 mm
- Y = clearance of bed = 96 in.
- Y_1 = clearance of wash troughs = 45 in.
- E_1 = expansion of anthracite [48] = 41 percent
- E_2 = expansion of sand [48] = 35 percent
- E_3 = expansion of garnet [48] = 10 percent
- G = minimum daily production = 1,000,000 gpd
- B = breakthrough index [2] = 1 (would equal 1×10^{-3} if d_e was expressed in cm)
- T = water temperature = 50°F
- U = deposit index [3] (equal to the value of the chi-square variable at the C/C_0 percentage point) = 22
- e_1 = average porosity for coal-sand-water mixture [8] = 0.85
- e_2 = average porosity for sand-garnet-water mixture [8] = 0.80
- D_1 = specific gravity of anthracite [7] = 1.5
- D_2 = specific gravity of sand [7] = 2.6
- D_3 = specific gravity of garnet [7] = 3.8
- D = specific gravity of water = 1.0

The media diameters will be established as constants to reduce the number of variables and constraints in the model. The use of these three constants eliminates three variables and eight constraints. To establish the diameters, the density of the anthracite-sand-water and sand-garnet-water mixtures must first be calculated per Equation 5.27.

$$D_{m1} = .85(1.0) + 1/2(1.5 + 2.6)(1 - .85) = 1.157$$

$$D_{m2} = .80(1.0) + 1/2(2.6 + 3.8)(1 - .80) = 1.44$$

The diameter ratio of adjacent media is calculated using Equation 5.28

$$\frac{1.6 d_1 (1.5 - 1.157)}{(2.6 - 1.157)} \leq d_2$$

$$.3803 d_1 \leq d_2 \quad (6.1)$$

$$\frac{1.6 d_2 (2.6 - 1.44)}{(3.8 - 1.44)} \leq d_3$$

$$.7864 d_2 \leq d_3 \quad (6.2)$$

If the diameter of garnet, d_3 , is set at 0.4mm, then d_1 and d_2 can be set per Equations 6.1 and 6.2 at 1.337 and .5086 respectively.

$$R = \frac{.05 (1 + .05/12)^{240}}{(1 + .05/12)^{240} - 1} = .0065701$$

and the backwash cost is

$$W = 15(6.7)(.00035) = \$.035175/\text{sq ft}$$

With the constants identified, step two is to determine the coefficients of the variables and set up the equations of the model.

Model for Sample Problem

The mathematical programming formulation of the reduced problem is given below:

Minimize

$$.00656 x_1 Q^{-1} + .002189 x_2 Q^{-1} + .03284 x_3 Q^{-1} + 2345 t^{-1} Q^{-1}$$

Subject to

$$142.86 Q^{-.34} t^{-1.7} d_e^{-.51} L^{.53} \leq 1$$

$$.00248 Q^{1.4} L^{.8} d_e^{-1.1} H^{-1} t \leq 1$$

$$H H_t^{-1} + 1.333 H_t^{-1} \leq 1$$

$$40 t^{-1} \leq 1$$

$$x_1 L^{-1} + x_2 L^{-1} + x_3 L^{-1} = 1$$

$$1.337 x_1 d_e^{-1} L^{-1} + .5086 x_2 d_e^{-1} L^{-1} + .4 x_3 d_e^{-1} L = 1$$

$$.125 H_t \leq 1$$

$$.03133 x_1 + .03x_2 + .02444x_3 \leq 1$$

$$3.0 Q^{-1} \leq 1$$

$$Q d_e^3 H_t L^{-1} \leq 1$$

The model now has nine variables, ten constraints, and 21 terms. As a geometric programming problem it thus has $21 - (9+1) = 11$ degrees of difficulty. This would be increased to 17 degrees if two inequalities were required for each of the two equality constraints. Two computer programs were tested on this problem but both experienced great difficulty in finding a solution. Computer programs for solving geometric programming

problems with large numbers of degrees of difficulty have as yet not been perfected [49].

To obtain an efficient solution (considering computational difficulty), the model was separated into two parts. The first part uses geometric programming to select the optimal flow rate, length of filter run, total depth of media, equivalent grain diameter, and head. The second part then uses the results from part one and a linear programming algorithm to compute the depth of anthracite, sand, and garnet. This approach makes use of the strong points of each technique. Successive application of more than one optimization technique in problem solution is becoming increasingly popular in solving complex problems.

Geometric Programming Solution

The geometric programming problem is

Minimize

$$2345 \, t^{-1} Q^{-1}$$

Subject to

$$142.86 \, Q^{-.34} t^{-1.7} d_e^{-.51} L^{.53} \leq 1$$

$$.00248 \, Q^{1.4} L^{.8} d_e^{-1.1} H^{-1} t \leq 1$$

$$7.5 \, H_t^{-1} \leq 1$$

$$.162 \, H \leq 1$$

$$40 \, t^{-1} \leq 1$$

$$Q \, d_e^3 \, H_t \, L^{-1} \leq 1$$

A lower bound of 7.5 feet and an upper bound of 6.17 feet were established on H_t and H in constraints three and four above in lieu of using the third and seventh constraints of the model. The lower bound on Q was also omitted. These actions avoided two terms and two degrees of difficulty. The problem above has seven terms, six variables, and zero degrees of difficulty. Application of the Edwards Geometric Programming Systems yields the following solution.

$$Q = 3.0 \text{ gpm/sq ft}$$

$$t = 40.0 \text{ time periods (10 hours)}$$

$$L = 25.0 \text{ inches}$$

$$d_e = 1.025 \text{ mm}$$

$$H = 6.17 \text{ ft}$$

$$H_t = 7.5 \text{ ft}$$

$$\text{Value of objective function} = \$18.9302/\text{million gal.}$$

Linear Programming Solution

The linear problem consists of the x_1 , x_2 , and x_3 terms from the model plus the following bounds on each media to ensure a proper mixture.

$$0.6L \leq x_1 \leq 0.75L$$

$$0.2L \leq x_2 \leq 0.3L$$

$$0.05L \leq x_3$$

The resulting linear programming problem is
Minimize

$$.002187 x_1 + .00073 x_2 + .01095 x_3$$

Subject to

$$1.337 x_1 + .5086 x_2 + .4 x_3 = 25.625$$

$$x_1 + x_2 + x_3 = 25$$

$$.03133 x_1 + .03 x_2 + .02444 x_3 \leq 1$$

$$x_1 \leq 18.75$$

$$x_1 \geq 15$$

$$x_2 \leq 7.5$$

$$x_2 \geq 5$$

$$x_3 \geq 1.25$$

The above problem was solved using MPS/360 yielding

$$x_1 = 15.8 \text{ in}$$

$$x_2 = 7.5 \text{ in}$$

$$x_3 = 1.7 \text{ in}$$

Value of objective function = \$.0586/million gallons

Summary of Results

The total of the objective function for the model is
 $18.9302 + .0586 = \$18.9888/\text{million gallons}$. The abscissa and ordinate
 values of the performance curves were calculated to test the

accuracy of the performance equations yielding

$$\frac{Q^{.2} d_e^{.3} t}{L^{.9}} = 2.77$$

$$\frac{U}{L} = 0.87$$

$$\frac{d_e^{1.4} (H_t - H_o)}{Q^{1.2} C_o^{.4} L^{1.7}} = .0022$$

The above results are very close to the values shown on the performance curves.

The influent concentration, C_o , and grain diameter, d_i , were varied for other runs of the model. Results of one of these runs (run 2) is compared in Table 6.1 with the previous data from run 1.

	Run 1	Run 2
	$C_0 = 20$	$C_0 = 5$
	$d_1 = 1.337$	$d_1 = 1.05$
	$d_2 = .5086$	$d_2 = .4$
	$d_3 = .4$	$d_3 = \text{none}$
Q	3.0	4.75
t	40.0	40.0
L	25.0	15.7
d_e	1.025	0.76
H	6.17	6.17
H_t	7.5	7.5
x_1	15.8	8.7
x_2	7.5	7.5
x_3	1.7	none
Obj. Function	18.9888	12.3409
$Q \cdot d_e^{.3} t/L^{.9}$	2.77	4.22
U/L	.87	1.78
$\frac{d_e^{1.4}(H_t - H_0)}{Q^{1.2} C_0^{.4} L^{1.7}}$.0022	.0032

Table 6.1 Output Comparisons

CHAPTER VII

CONCLUSIONS AND RECOMMENDATIONS

Filter Cost Model

Mathematical cost modeling is a practical approach to filter design. Such models are of course dependent on the availability of curves or equations that can predict filter performance. Hsiung, as indicated previously, has provided a method for developing the needed curves into a cost model. The cost model and optimization solution methodology presented in this report will also be useful in the determination of the optimum operational policies for an existing filtration plant.

Refinements to the model are encouraged. Filtration is a highly complex process and there are no other items and constraints that could be added to the model to make it more universal. This report merely introduces the technique in the hope that experienced environmental design engineers can improve on its form and find areas where it can be practically applied.

Geometric Programming Optimization

It is reasonable to expect the filter performance equations to be as complex as the process they are describing. Without geometric programming, linear or quadratic approximations to the equations would have to be made before a solution can be obtained. Geometric programming, however, can be used to solve these highly nonlinear equations and provide accurate answers. Also, very little additional computations are

necessary to determine the new value of the objective function whenever prices or cost coefficients change.

Geometric programming would be even more valuable if a general computer algorithm can be developed to solve problems with many degrees of difficulty. The Edwards GPS is a good basis for such a program, but it sorely needs a simplex subroutine to find an initial feasible solution in the linear dual space. Without this initial point, it experiences difficulty at times locating the feasible region. There are several search routines that could be used for the degree of difficulty problems, but one must be careful in compromising between efficiency and level of significance. General search programs with many significant figures are desirable, but such programs can consume a great deal of computer time.

Combining Optimization Techniques

Arguments might be offered that the separation of the problem into two subproblems does not guarantee a global optimum. These arguments probably have some basis, but the global optimum is an aspiration that few design engineers actually achieve. In modeling, it is usually necessary to loop between deductive tractability and assumptions for characterization.

The combination of operations research techniques can have wide application and be extremely useful in the solution of real world problems. Combining geometric programming with linear programming, dynamic programming, search, simulation, etc. can easily lead to the solution of many complex problems.

Future Extensions

The model developed in this research needs to be tested with actual data and experiments run to compare predictions with experimental results. A sensitivity analysis also needs to be performed to measure the effect of each variable and constant on the final results.

The modeling technique itself can be used to predict performance of wastewater and diatomaceous earth filters. Filtration theory for diatomaceous earth filters is much more refined than that for sand or multimedia filters, and a generally accepted set of performance equations are available [14]. The model can also be extended in water filtration to include more of the costs involved in the construction of a new filtration plant.

The geometric programming technique has a wide range of potential applications in environmental engineering including hydrology, chemical treatment processes, water resource analysis, and others.

Mathematical modeling is an inexpensive method of describing and testing a system. It can help the engineer to build conceptual structures out of perceptual confusion.

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APPENDIX A
PROGRAM LISTING


```

C.....SIG(M), M = 2,...CAPM, (=+ OR -1) = UPPER BOUND ON CONSTRAINT M.
C.....
C * * * E N T E R I N G   F L O W C H A R T   B L O C K   A   * * * *
100 READ (5,1) CAPM,CAPN,(TSSUBM(M),M=1,CAPM),(SIG(M),M=1,CAPM)
1   FORMAT (22I2)
C.....
C.....   CLOCK IS NOT AVAILABLE IN WATFOR.
C.....
C.....   92 CALL CLOCK(HRS,AMIN,ASEC)
C * * * E N T E R I N G   F L O W C H A R T   B L O C K   B   * * * *
WRITE(6,30)
30 FORMAT('1.T43,***POLYNOMIAL MINIMIZATION BY GEOMETRIC PROGRAMMING
1***./)
C.....
C.....   READ-IN / READ-OUT TITLE.
C.....
C.....   READ (5,102) (NOTE(I),I=1,10)
102 FORMAT(10A8)
WRITE(6,110) (NOTE(I),I=1,10)
110 FORMAT(T43,10A8)
C.....
C.....   READ LOWER LIMIT FOR STEP SIZE IN LETTO PATTERN SEARCH
C.....   SUBROUTINE.   DEFAULT OPTION/XLLIM = 10.**(-6)
C.....   READ LPRINT, AN INTEGER FROM 0 TO 4.   HIGHER VALUES CAUSE
C.....   LESS INFORMATION ON THE PATTERN SEARCH TO BE PRINTED OUT.
C.....
READ(5,210) XLLIM,LPRINT
210 FORMAT(E13.0,I1)
IF(XLLIM.LE.0.0) XLLIM = 10.**(-6)
I = CAPM-1

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00089

      WRITE (6,103)CAPN,I
103  FORMAT(///T40,'PRIMAL PROGRAM.'/T40,15('='))///
      1T40,'THE NUMBER OF PRIMAL VARIABLES = 'I2/
      1T40,'THE NUMBER OF PRIMAL CONSTRAINTS = 'I2)
      IF(CAPN.LE.0.OR.CAPM.LE.0) CALL BLOWUP(1,&104)
      IF(CAPN.GT.10.OR.CAPM.GT.10) CALL BLOWUP(2,&104)
      WRITE(6,35)
35  FORMAT(///T55,'*** GENERALIZED POLYNOMIAL ***'///T40,'ROW'
      1T46,'TERM'T54,'COEFFICIENT'T69,'PRIMAL'T81,'EXPONENT'T94,'SIGNUM'/'
      1T56,'OF TERM'T68,'VARIABLE'T80,'OF VARIABLE'T94,'OF ROW'/'
      1 T40,60('.'/)/T40,'MINIMIZEO'//)
      CAPT = 0
      CAPNP1 = CAPN+1
      CAPNP2 = CAPN+2
      DO 8 M=1,CAPM
      IF(IABS(SIG(M)).NE.1.OR.TSUBM(M).LE.0) CALL BLOWUP(1,&104)
      IF(TSUBM(M).GT.10) CALL BLOWUP(2,&104)
      I =TSUBM(M)
      DO 4 T=1,I
      CAPT = CAPT+1
      C..... READ XCON, THE COEFFICIENT OF THE T-TH TERM AND THE M-TH ROW.
      C.....
      C.....
      READ(5,2) XCON(M,T)
      2  FORMAT(E25.0)
      WRITE (6,36) M,T,XCON(M,T)
      36  FORMAT(/T40,I2,T47,I2,T51,G15.8)
      IF(XCON(M,T).EQ.0.0) CALL BLOWUP(1,&104)
      XABS =DABS(XCON(M,T))
      SIGMA(M,T) = XCON(M,T)/XABS

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XCON(M,T) = XABS

C..... READ A(M,T,N), THE EXPONENT OF THE N-TH VARIABLE IN THE T-TH
C..... TERM OF THE M-TH ROW OF THE GENERALIZED POLYNOMIAL. IF LAST=1,
C..... N IS CONSIDERED TO BE THE LAST VARIABLE IN THAT TERM.
C.....
C.....
5 READ(5,3) N,A(M,T,N),LAST
3 FORMAT (I2,E16.0,I2)
WRITE (6,37) N,A(M,T,N)
37 FORMAT(' ',T67,'* X(',I2,') ** ',G15.8/)
IF(N.LE.0.DR.N.GT.CAPN.DR.A(M,T,N).EQ.0) CALL BLOWUP(1,E104)
IF(LAST.EQ.0) GO TO 5
IF(T.NE.1) WRITE(6,301)
301 FORMAT(T62,'.PLUS. '/')
4 CONTINUE
IF(M.EQ.1.AND.SIG(1).GT.0) WRITE(6,32)
32 FORMAT(T40,'SIGN OF PRIMAL FUNCTION AT MINIMUM ASSUMED TO BE POSIT
1IVE',/T40,60(' '))
IF(M.EQ.1.AND.SIG(1).LT.0) WRITE(6,132)
132 FORMAT(T40,'SIGN OF PRIMAL FUNCTION AT MINIMUM ASSUMED TO BE NEGAT
1IVE',/T40,60(' '))
IF(M.EQ.1.AND.CAPM.GT.1) WRITE (6,31)
31 FORMAT(T40,'SUBJECT TO')
IF(M.GT.1) WRITE(6,33) SIG(M)
33 FORMAT(' ',T92,'.LE.',T98,I2/T40,60(' '))
8 CONTINUE

C..... THE NUMBER OF DEGREES OF DIFFICULTY IS THE NUMBER OF TERMS IN
C..... THE OBJECTIVE FUNCTION AND ALL CONSTRAINTS MINUS THE NUMBER
C..... OF PRIMAL VARIABLES MINUS ONE, AND MAY NOT BE LESS THAN ZERO

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C..... OR GREATER THAN TWENTY.
C.....
  IDOIF = CAPT-CAPN-1
  WRITE(6,34) CAPT,IDOIF
34 FORMAT(///T40,'THE TOTAL NUMBER OF TERMS = 'I2/
1 T40,'THE NUMBER OF 'DEGREES OF DIFFICULTY' = 'I2)
  IF(IDOIF.LT.0.OR.IDOIF.GT.20) CALL BLOWUP(3,6104)

C..... BEGIN CONSTRUCTION OF MATRIX OF COEFFICIENTS OF DUAL
C..... VARIABLES.
C.....
C * * * E N T E R I N G F L O W C H A R T B L O C K C * * * *
CAPT1 = CAPT+1
94 IT=TSUBM(1)
DO 6 T=1,IT
6 B(1,T)=SIGMA(1,T)
  IT = IT + 1
DO 316 T = IT,CAPT1
316 B(1,T) = 0.0
  TT = 0
DO 7 M=1,CAPM
  IT = TSUBM(M)
DO 7 T=1,IT
  TT = TT+1
DO 7 N=1,CAPN
7 B(N+1,TT) = SIGMA(M,T)*A(M,T,N)
DO 290 I = 2,CAPNP1
290 B(1,CAPT1) = 0.0
  B(1,CAPT1) = SIG(1)
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C..... MATRIX OF COEFFICIENTS OF DUAL VARIABLES IS COMPLETED.
C..... CALL MAT TO ARRANGE IDENTITY MATRIX IN LEFT SQUARE MATRIX OF B.
C.....
C * * * * * E N T E R I N G   F L O W C H A R T   B L O C K   D   * * * * *
C * * * * * CALL MAT(CAPNP1,CAPT1,£104)
C.....
C..... IF ALL VALUES ON THE RIGHT HAND SIDE ARE NEGATIVE, SIG(1) IS
C..... OF THE WRONG SIGN AND THE MATRIX MUST BE RE-SOLVED WITH
C..... SIG(1) = -SIG(1).
C.....
C..... DO 93 I=1,CAPNP1
C..... IF(B(ROW(I),COL(CAPT1))) 93,93,96
C.....
C..... 93 CONTINUE
C..... 193 DO 95 I = 1,32
C..... ROW(I) = I
C..... 95 COL(I) = I
C..... SIG(1) = -SIG(1)
C..... WRITE (6,220)
C..... 220 FORMAT('1./T40,'SIGN OF DUAL FUNCTION AT MAXIMUM IS NOT THE SAME A
C..... IS SPECIFIED FOR PRIMAL.'/T40,'SIG(1) = -SIG(1) ',/T40,'EXECUTION CO
C..... INTINUES')
C..... GO TO 94
C..... 96 DO 500 I = 1,CAPNP1
C..... DO 500 J = 1,CAPT1
C..... 500 B(I+11,J) = B(ROW(I),COL(J))
C..... DO 501 I = 1,CAPNP1
C..... DO 501 J = 1,CAPT1
C..... 501 B(I,J) = B(I+11,J)
C * * * * * E N T E R I N G   F L O W C H A R T   B L O C K   E   * * * * *
C * * * * * WRITE (6,40) CAPT,CAPNP1

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00180 40 FORMAT('1'//T40,'DUAL PROGRAM.'/T40,13('=')/T40 ,
00181 1'THE NUMBER OF DUAL VARIABLES = 'I2/
00182 1 T40,'THE NUMBER OF DUAL EQUATIONS = 'I2////
00183 2 T42,'*** EXTENDED MATRIX OF COEFFICIENTS OF DUAL VARIABLES ***'/
00184 1 T40,60('-.'))
00185 IF(CAPT1.GT.12) GO TO 303
00186 DO 25 M=1,CAPNP1
00187 25 WRITE(6,41)(B(M,I),I=1,CAPT1)
00188 41 FORMAT(/T40,13F7.3/T43,9F7.3)
00189 GO TO 302
00190 303 DO 304 M = 1,CAPNP1
00191 304 WRITE(6,305)(B(M,I),I=1,CAPT1)
00192 305 FORMAT(/T15,16F7.3/T18,16F7.3)
00193 C..... IF DEGREES OF DIFFICULTY EXIST, PATTERN SEARCH IS CALLED WITH
00194 C..... IF DEGREES OF DIFFICULTY EXIST, PATTERN SEARCH IS CALLED WITH
00195 C..... IF DEGREES OF DIFFICULTY EXIST, PATTERN SEARCH IS CALLED WITH
00196 C..... IF DEGREES OF DIFFICULTY EXIST, PATTERN SEARCH IS CALLED WITH
00197 302 MPT = CAPM+CAPT
00198 NCON = MPT-1
00199 IF(1DDIF.GT.0) WRITE(6,42) 1DDIF,NCON,XLLIM,LPRINT
00200 42 FORMAT(T40,60('-.'))//1'T40,'PATTERN SEARCH.'/T40,15('=')////
00201 1T40,'THE NUMBER OF VARIABLES = ' I2/
00202 1 T40 ,'THE NUMBER OF CONSTRAINTS = 'I2/
00203 2 T40,'THE LOWER LIMIT ON STEP SIZE = ',1PE12.4/
00204 1 T40,'PRINT OPTION = 'I2////
00205 C * * * E N T E R I N G F L O W C H A R T B L O C K F * * * *
00206 IF(1DDIF.GT.0) CALL LETTO(6104)
00207 C * * * E N T E R I N G F L O W C H A R T B L O C K G * * * *
00208 CALL VAL(X,DUAL)
00209 C * * * E N T E R I N G F L O W C H A R T B L O C K H * * * *

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00210 WRITE(6,60) DUAL
00211 60 FORMAT(T40,60('---'))//1,T40,'DUAL SOLUTION.'/T40,14('---')////
00212 1T40,'THE VALUE OF THE DUAL FUNCTION AT A MAXIMUM POINT = '/T80,G16
00213 1.8//T40,'THE CORRESPONDING DUAL VARIABLES AREO'//T40,60('---')//
00214 1T60,'DUAL',
00215 T78,'VALUE'/T58,'VARIABLE'/T40,60('---')
00216 J = CAPT+CAPM
00217 WRITE(6,179)(I,W(I),I=1,J)
00218 179 FORMAT(T60,'W('',I2,'') = ',T70,G16.8)
00219
00220 C..... BEGIN CONSTRUCTION OF MATRIX OF EXPONENTS OF PRIMAL VARIABLES.
00221 C.....
00222 C * * * ENTERING FLOWCHART BLOCK I * * * *
00223 IT = TSUBM(1)
00224 DO 50 T=1,IT
00225 ROW(T) = T
00226 COL(T) = T
00227
00228 C..... A DUAL VARIABLE CLOSE TO ZERO CAUSES EXCESSIVE ROUND-OFF ERROR.
00229 C..... THAT ROW OF THE MATRIX IS SET TO ZERO, AND WILL BE DISCARDED
00230 C..... BY MAT.
00231 C.....
00232 IF(W(T).LT.10.**(-30)) GO TO 200
00233
00234 C..... IF SIG(1)*DUAL IS NEGATIVE, LET SIG(1) = -SIG(1) AND GO BACK
00235 C..... TO THE DUAL SPACE. HOWEVER, IF SIG(1)*DJAL = 0, DISREGARD.
00236 C.....
00237 IF(SIG(1)*DUAL) 193,200,230
00238 DO 51 N = 1,CAPN
00239 B(T,N) = A(1,T,N)

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THE EQUATIONS FOR SOLUTION OF THE PRIMAL VARIABLES

C..... ARE LINEAR IN LN(X)
C.....
C..... B(IT,CAPNP1) = DLOG(SIG(1)*DUAL*W(T)/XCON(1,T))

GO TO 49
200 DO 201 N = 1,CAPNP1
201 B(T,N) = 0.0
49 CONTINUE
50 CONTINUE

IF(CAPM.EQ.1) GO TO 54

DO 52 M=2,CAPM

I=TSUBM(M)

DO 52 T = 1,I

IT = IT+1

RDW(IT) = IT

COL(IT) = IT

ATEMP = W(CAPT+M)*XCON(M,T)

IF(DABS(ATEMP).LT.10.**(-30)) GO TO 202

ATEMP = W(IT)/ATEMP

IF(ATEMP.LT.10.**(-30)) GO TO 202

B(IT,CAPNP1) = DLOG(ATEMP)

DO 53 N = 1,CAPN

53 B(IT,N) = A(M,T,N)

GO TO 59

202 DO 203 N = 1,CAPNP1

203 B(IT,N) = 0.0

59 CONTINUE

52 CONTINUE

C..... MATRIX OF COEFFICIENTS OF LOGS OF PRIMAL VARIABLES IS COMPLETE.
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C..... CALL MAT TO SOLVE SIMULTANEOUS EQUATIONS FOR PRIMAL VARIABLES.
C.....
C*** * ENTERING FLOWCHART BLOCK J * * * * *
54 CALL MAT(CAPT,CAPNP1,£104)
   DO 80 I = 1,CAPN
   IF(B(ROW(I),COL(CAPNP1)),GT.174.673) B(ROW(I),COL(CAPNP1))=174.673
80 B(ROW(I),COL(CAPNP1)) =DEXP(B(ROW(I),COL(CAPNP1)))
   WRITE (6,74)
C*** * ENTERING FLOWCHART BLOCK K * * * * *
74 FORMAT(T40,60('---'))///T40,'PRIMAL SOLUTION.'/T40,16('---'))///
   1T40,'THE VALUES OF THE PRIMAL VARIABLES AT A MINIMUM POINT ARE ' /
   1T40,60('---'))/T60,'PRIMAL',T75,'VALUE'/T58,'VARIABLE'/T40,60('---'))
   WRITE (6,73) (I,B(ROW(I),COL(CAPNP1)),I=1,CAPN)
73 FORMAT(T60,'X(',I2,') = ',T70,G16.8)
   WRITE (6,77)
77 FORMAT ( T40,60('---'))///T40 ,
   1,'THE CONTRIBUTION OF EACH TERM TO THE PRIMAL FUNCTION ISO' //
   1 T40,60('---'))/T60,'TERM',T75,'CONTRIBUTION'/T40,60('---'))

C..... SUBSTITUTE PRIMAL VARIABLES INTO GENERALIZED POLYNOMIAL.
C.....
C.....
ITT = 0
DO 71 M = 1,CAPM
Y = 0.0
IT = TSUBM(M)
DO 79 T=1,IT
ITT = ITT + 1
YTERM = 1.0
DO 7C N = 1,CAPN
C.....

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00300      C..... IF C3EFFICIENT A = 0.0, X**A = 1.0, UNLESS X.LT.0,
00301      C..... IN WHICH CASE THE TERM BLOWS UP.
00302      C.....
00303      C..... IF (A(M,T,N)) 260,70,170
00304      260 IF (B(ROW(N),COL(CAPNP1)))170,261,170
00305      261 YTERM = YTERM*10.**{(30)
00306      GO TO 70
00307      170 YTERM = YTERM*B(ROW(N),COL(CAPNP1))*A(M,T,N)
00308      70 CONTINUE
00309      YTERM = SIGMA(M,T)*YTERM*XCON(M,T)
00310      WRITE (6,178) ITT,YTERM
00311      178 FORMAT('62,I2,T70,G16.8')
00312      79 Y = Y+YTERM
00313
00314      C..... YY = VALUE OF THE PRIMAL OBJECTIVE FUNCTION.
00315      C.....
00316      C..... IF (M.EQ.1) YY = Y
00317      71 WRITE (6,76) M,Y
00318      76 FORMAT('T40,' THE VALUE OF ROW 'I2,' = 'G16.8, /)
00319      WRITE (6,75) YY
00320      75 FORMAT('T40,60('---')////T40,' THE VALUE OF THE PRIMAL FUNCTION = '
00321      1 G16.8/)
00322
00323      C..... COMPUTE ABSOLUTE AND RELATIVE ERROR BETWEEN PRIMAL AND DUAL
00324      C..... FUNCTIONS. IN THE CONVEX CASE, THE SOLUTION IS GUARANTEED
00325      C..... TO LIE BETWEEN THESE FUNCTIONS.
00326      C.....
00327      D = DABS(YY-DUAL)
00328      E = (YY+ DUAL)/2.0
00329      F = D/E

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WRITE(6,78) DUAL,D,D,E,F
78 FORMAT(T40,'THE VALUE OF THE DUAL FUNCTION = 'G16.8//
1T40,'THE ABSOLUTE VALUE OF THE DIFFERENCE = 'G16.8//
1T40,'THE RELATIVE DIFFERENCE = 'G13.6,'/'G13.6//T64,'= '2PF10.6,
1' PER CENT'//)
DO 310 M = 1,CAPM
IF(SIG(M)) 315,311,311
311 IT = TSUBM(M)
DO 310 T = 1,IT
IF(SIGMA(M,T)) 315,310,310
310 CONTINUE
WRITE(6,313)
GO TO 104
315 WRITE(6,312)
313 FORMAT(/T40,'PRIMAL PROBLEM IS CONVEX. GLOBAL MINIMUM GUARANTEED.
1'//)
312 FORMAT(/T40,'PRIMAL PROBLEM IS NON-CONVEX.'
1 /T40,'GLOBAL MINIMUM IS NOT GUARANTEED.'//T40,
1 'THIS STATIONARY POINT MAY BE A LOCAL MINIMUM, SADDLE POINT, OR M
1AXIMUM.'//)
C.....
C..... TO TERMINATE, PUT 1 IN COL 2 AND LEAVE REST BLANK.
C..... CAPM MUST NOT = 0 OR ERROR WILL RESULT.
C..... IF CAPN = 0, PROGRAM WILL TERMINATE.
C..... OTHERWISE, IT WILL GO BACK TO 92.
C.....
104 READ (5,1) CAPM,CAPN,(TSUBM(M),M=1,CAPM),(SIG(M),M=1,CAPM)
C.....
C..... CLOCK IS NOT AVAILABLE IN WATFOR
C.....

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00360 CALL CLOCK(HRS,AMIN,BSEC)
00361 ATIME = BSEC - ASEC
00362 WRITE(6,111) ATIME
00363 111 FORMAT('43',ELAPSED TIME = ',F6.3,' SECONDS.'/1')
00364 IF(CAPN) 101,101,250
00365
00366 C..... RE-INITIALIZE CERTAIN MATRICES TO 0.
00367 C..... RE-INITIALIZE CERTAIN MATRICES TO 0.
00368 C..... CLEAR AND CLEAR ARE NOT AVAILABLE IN WATFOR
00369 C.....
00370 250 CALL CLEAR(A(1,1,1),1000)
00371 CALL CLEAR(XCON(1,1),100)
00372 CALL CLEAR(SIGMA(1,1),100)
00373 DO 91 I = 1,32
00374 ROW(I) = 1
00375 COL(I) = 1
00376 91 CONTINUE
00377 GO TO 92
00378 101 STOP
00379 END
00380 SUBROUTINE MAT(NROWS,NCOLS,*)
00381 C..... THIS SUBROUTINE PLACES AN IDENTITY MATRIX IN THE LEFT-HAND
00382 C..... SQUARE MATRIX OF B , WHICH MAY BE SQUARE OR RECTANGULAR.
00383 C.....
00384 C..... IMPLICIT REAL*8(A-8,D-H,O-Q,U-Z), INTEGER(C,R-T)
00385 COMMON/TOM/A(10,10,10),B(31,32),XCON(10,10),SIGMA(10,10),
00386 1ROW(32),COL(32),TSUBM(10),CAPM,CAPN,CAPNP1,CAPNP2,CAPT,CAPTPI
00387 NCM1 = NCOLS - 1
00388 C.....
00389 C..... CREATE UPPER-TRIANGULAR MATRIX.

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00390 C.....
00391 C** * E N T E R I N G   F L O W C H A R T   B L O C K   L   * * * * *
00392 C.....
00393 C..... DO 10 IPIV = 1,NROWS
00394 C.....
00395 C..... DIVIDE ROW BY PIVOT ELEMENT.
00396 C.....
00397 C..... IPIV = IPIV+1
00398 C..... IF(B(ROW(IPIV),COL(IPIV)).EQ.1.0) GO TO 23
00399 C..... DO 31 TT=IPIV,NCM1
00400 C..... DO 11 T = IPIV,NROWS
00401 C.....
00402 C..... IF PIVOT ELEMENT = 0.0, EXCHANGE ROWS.
00403 C.....
00404 C..... IF PIVOT ELEMENT = 0.0, EXCHANGE ROWS.
00405 C..... IF PIVOT ELEMENT = 0.0, EXCHANGE ROWS.
00406 C..... IF PIVOT ELEMENT = 0.0, EXCHANGE ROWS.
00407 C..... IF ALL ROWS HAVE BEEN EXCHANGED, TRY EXCHANGING COLUMNS.
00408 C.....
00409 C..... IF ALL ROWS HAVE BEEN EXCHANGED, TRY EXCHANGING COLUMNS.
00410 C..... IF ALL ROWS HAVE BEEN EXCHANGED, TRY EXCHANGING COLUMNS.
00411 C..... IF ALL ROWS HAVE BEEN EXCHANGED, TRY EXCHANGING COLUMNS.
00412 C..... IF ALL ROWS HAVE BEEN EXCHANGED, TRY EXCHANGING COLUMNS.
00413 C..... IF ALL ROWS HAVE BEEN EXCHANGED, TRY EXCHANGING COLUMNS.
00414 C..... IF ALL ROWS HAVE BEEN EXCHANGED, TRY EXCHANGING COLUMNS.
00415 C..... IF ALL ROWS HAVE BEEN EXCHANGED, TRY EXCHANGING COLUMNS.
00416 C..... IF ALL ROWS HAVE BEEN EXCHANGED, TRY EXCHANGING COLUMNS.
00417 C..... IF ALL ROWS HAVE BEEN EXCHANGED, TRY EXCHANGING COLUMNS.
00418 C..... IF ALL ROWS HAVE BEEN EXCHANGED, TRY EXCHANGING COLUMNS.
00419 C..... IF ALL ROWS HAVE BEEN EXCHANGED, TRY EXCHANGING COLUMNS.

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00420      IT = COL(IPIV)
00421      COL(IPIV) = COL(TT+1)
00422      31 COL(TT+1) = IT
00423      C * * * * * E N T E R I N G   F L O W C H A R T   B L O C K   R   * * * * *
00424
00425      12 DO 13 T = IPPI, NCOLS
00426      13 B(ROW(IPIV),COL(T))=B(ROW(IPIV),COL(T))/B(ROW(IPIV),COL(IPIV))
00427      B(RCW(IPIV),COL(IPIV)) = 1.0
00428
00429      C..... THE ROW HAS BEEN DIVIDED BY THE PIVOT ELEMENT.
00430      C..... IF THIS COMPLETES THE UPPER-TRIANGULAR MATRIX , GO TO 14.
00431      C.....
00432      23 IF(IPIV.EQ.NROWS.OR.IPIV.EQ.NCMI) GO TO 14
00433
00434      C..... PLACE ZEROS IN THE ELEMENTS BELOW THE PIVOT ELEMENT.
00435      C.....
00436      DO 21 N = IPPI,NROWS
00437
00438      C..... IF THE ELEMENT IS ALREADY ZERO, SKIP THAT ROW.
00439      C.....
00440      IF(B(ROW(N),COL(IPIV))) 41,21,41
00441      C * * * * * E N T E R I N G   F L O W C H A R T   B L O C K   S   * * * * *
00442      41 IT = NCOLS+1-IPIV
00443      DO 15 I=1,IT
00444      T = NCOLS+1-I
00445      15 B(ROW(N),COL(T))=B(ROW(N),COL(T))/B(ROW(N),COL(IPIV))
00446      1 -B(RCW(IPIV),COL(T))
00447      21 CONTINUE
00448      C * * * * * T H I S   I S   F L O W C H A R T   P O I N T   I O   * * * * *
00449      10 CONTINUE
00450      C.....

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C..... UPPER-TRIANGULAR MATRIX COMPLETED.
C..... IF THERE IS ONLY ONE ROW OR ONE COLUMN, RETURN.
C.....
C..... 14 IF(NROWS.EQ.1.OR.NCM1.EQ.1) RETURN
C.....
C..... IF THERE ARE MORE ROWS THAN COLUMNS, DISCARD THE BOTTOM ROWS.
C.....
C..... IF(NROWS.GT.NCM1) NROWS = NCM1
C..... NRPI = NRJWS + 1
C.....
C..... PLACE ZEROS IN UPPER ELEMENTS.
C..... PIVOT ELEMENTS START AT THE BOTTOM.
C.....
C..... ENTERING FLOWCHART BLOCK T * * * * *
C * * * DO 20 T = 2, NRJWS
C * * * IPIV = NRJWS+2-T
C * * * IPM1 = IPIV-1
C * * * DO 20 N=1, IPM1
C..... IF THERE IS ALREADY A ZERO IN THIS ELEMENT, SKIP THIS ROW.
C.....
C..... IF(B(ROW(N),COL(IPIV))) 42,20,42
C..... IF(B(ROW(N),COL(IPIV))) 42,20,42
C * * * ENTERING FLOWCHART BLOCK U * * * * *
C * * * 42 BCOEF = B(ROW(N),COL(IPIV))/B(ROW(IPIV),COL(IPIV))
C * * * DO 22 TCOL=NRPI,NCOLS
C * * * B(ROW(N),COL(TCOL)) = B(ROW(N),COL(TCOL))
C * * * 1 - B(ROW(IPIV),COL(TCOL))*BCOEF
C * * * 22 CONTINUE
C * * * B(ROW(N),COL(IPIV)) = 0.0
C * * * 20 CONTINUE

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00480 C..... IDENTITY MATRIX COMPLETED.
00481 C.....
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00508 C.....
00509 C.....

104 WRITE (6,45)
45 FORMAT(/T40,'SINGULARITY DISCOVERED IN MATRIX*****'//)
DO 44 I = 1,NROWS
44 WRITE (6,43) (B(ROW(I),COL(J)),J = 1,NCOLS)
43 FORMAT(T43,13F7.3/T43,9F7.3)
WRITE(6,46)
46 FORMAT(/T40,'THIS IS A PATHOLOGICAL PROBLEM.'//
1 T40,'IT IS SUGGESTED THAT THE PRIMAL PROBLEM EITHER HAS NO
1 /T40,'MINIMUM OR HAS AN OBVIOUS AND 'TRIVIAL' MINIMUM.'/
1 T40,'EXECUTION TERMINATED*****'//)
RETURN 1
END
SUBROUTINE BLOWUP(I,*)
C.....
C..... THIS SUBROUTINE IS CALLED FROM THE MAIN PROGRAM AND WILL
C..... TERMINATE PROBLEM EXECUTION.
C.....
C.....
IMPLICIT REAL*8(A-B,D-H,O-Q,U-Z), INTEGER(C,R-T)
GO TO (1,2,3),I
1 WRITE(6,10)
GO TO 100
2 WRITE(6,20)
GO TO 100
3 WRITE(6,30)
10 FORMAT(/T43,'ABOVE DATA IMPLIES DATA CARD ERROR.'//)
20 FORMAT(/T43,'ABOVE DATA EXCEEDS UPPER LIMITS ON SIZE OF PROBLEM')

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D = 0.01
ADAF = 1.0
CFSC = 0.1314159
CASS = CFSC
ISTOP = 2
IFEAS = 1
KC = 0
LACT = 0
CALL SUMS(X, ZBASE)
IF(LPRINT.LE.3) WRITE (6, 1110)
IF(LPRINT.LE.3) INITIAL POINT.
1110 FORMAT( T40, 'INITIAL POINT.
1110 IF(LPRINT.LE.3) CALL OUTPUT (X, ZBASE)
ZTEMP = ZBASE
DO 110 I = 1, N
XTEMP(I) = X(I)
110 XBASE(I) = X(I)
IF (IFEAS .EQ. 1) GO TO 111
IF(LPRINT.LE.2) WRITE (6, 5110)
5110 FORMAT(T40, 'INITIAL POINT NOT IN FEASIBLE REGION.
CALL SUMC(X, CBASE)
CTEMP = CBASE
CFSCI = CFSC
C*** THIS IS POINT A WHICH APPEARS IN FLOWCHART, FIG. 8.
5111 DO 5112 I = 1, N
5112 XTEST(I) = XBASE(I)
5114 DO 5115 I = 1, N
XTEST(I) = XTEMP(I) + CFSCI
CALL SUMC (XTEST, CTEST)
IF (CTEST .GT. CTEMP) GO TO 5135
XTEST(I) = XTEMP(I) - CFSCI

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5174 XBASE(I) = XTEST(I)
      GO TO 5175
C*** THIS IS POINT B WHICH APPEARS IN FLOWCHART, FIG. 10.
5146 DO 5147 I = 1,N
      X(I) = XTEMP(I)
5147 XBASE(I) = XTEMP(I)
C*** THIS IS POINT F WHICH APPEARS IN FLOWCHART, FIG. 10.
5176 IF(LPRINT.LE.2) WRITE(6,5177)
5177 FORMAT(T40,'FEASIBLE POINT HAS BEEN FOUND.')
```

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      CALL VAL(X, ZBASE)
      ZTEMP = ZBASE
      GO TO 111
5210 IF(LPRINT.LE.1) WRITE(6,5211)
5211 FORMAT(T40,'MORE FEASIBLE POINT NOT FOUND IN LOCAL SEARCH.')
```

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      IT40,'REDUCE STEP SIZE (CFSCI).')
      CTEMP = CBASE
      DO 5216 I = 1,N
5216 XTEMP(I) = XBASE(I)
      CFSCI = CFSCI/3.0
      IF(CFSCI.GE.XLLIM) GO TO 5111
      IF(LPRINT.LE.3) WRITE(6,5270)
5270 FORMAT(T40,'NO FEASIBLE POINT FOUND. PROBLEM TERMINATED.')
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      RETURN
C*** THIS IS POINT D WHICH APPEARS IN FLOWCHART, FIG. 10.
5172 IF(LPRINT.LE.1)WRITE(6,6172)
6172 FORMAT(T40,'BETTER POINT NOT FOUND IN PATTERN MOVE.')
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      IT40,'RETURN TO BEST POINT IN LOCAL SEARCH.')
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      CBASE = CTEMP
      DO 5195 I = 1,N
      X(I) = XTEMP(I)

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5195 XBASE(I) = XTEMP(I)
GO TO 5111
C*** THIS IS PJINT 111 WHICH APPEARS IN FLOWCHART, FIG. 11.
111 DO 112 I = 1, N
XTEMP(I) = XBASE(I)
112 XTEST(I) = XBASE(I)
ZTEMP = ZBASE
114 DO 115 I = 1, N
XTEST(I) = XTEMP(I) + CFSC
CALL SUMS(XTEST, ZTEST)
IF (ZTEST .GT. ZTEMP) GO TO 135
XTEST(I) = XTEMP(I) - CFSC
CALL SUMS (XTEST, ZTEST)
IF (ZTEST .LE. ZTEMP) GO TO 115
135 ZTEMP = ZTEST
XTEMP(I) = XTEST(I)
115 CONTINUE
C*** THIS IS PJINT 6 WHICH APPEARS IN FLOWCHART, FIG. 12.
CALL SUMS(XTEMP, ZTEST)
IF (LPRINT.LE.0) CALL OUTPUT (XTEMP, ZTEST)
145 IF (IFEAS.EQ. 2) GO TO 146
150 IF (ZTEST.LE. ZBASE + XLLIM) GO TO 149
ZTEMP = ZTEST
CASS = CFSC
ISTOP = 2
ADAF = 1
DO 160 I = 1, N
160 XTEST(I) = XTEMP(I)*2. -X(I)
CALL SUMS (XTEST, ZTEST)
IF (LPRINT.LE.0) CALL OUTPUT (XTEST,ZTEST)

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170 IF (ZTEST.LE. ZTEMP+ XLLIM) GO TO 172
171 IF (IFEAS.EQ. 2) GO TO 173
  ZBASE = ZTEST
  ZTEMP = ZTEST
  ADAF = 1.0
  DO 180 I = 1, N
    X(I) = XTEMP(I)
180 XBASE(I) = XTEST(I)
    IF (LPRINT.LE.1) WRITE (6,1180)
1180 FORMAT(T40, 'LET PATTERN MOVE STAND.')
    GO TO 111
173 IF (LPRINT.LE.1) WRITE (6, 1173)
1173 FORMAT(T40, 'CONSTRAINT CONTACTED IN PATTERN MOVE.', /
  T40, 'RETURN TO BEST POINT IN LOCAL SEARCH.')
  GO TO 190
C*** THIS IS POINT K WHICH APPEARS IN FLOWCHART, FIG. 13.
172 IF (LPRINT.LE.1) WRITE (6, 1172)
1172 FORMAT(T40, 'BETTER POINT NOT FOUND IN PATTERN MOVE.' /
  T40, 'RETURN TO BEST POINT IN LOCAL SEARCH.')
C*** THIS IS POINT L WHICH APPEARS IN FLOWCHART, FIG. 13.
190 ZBASE = ZTEMP
  DO 195 I = 1, N
    X(I) = XTEMP(I)
195 XBASE(I) = XTEMP(I)
  GO TO 111
151 IF (LPRINT.LE.1) WRITE (6, 1151)
1151 FORMAT(T40, 'MULTIPLE CONSTRAINTS CONTACTED IN LOCAL SEARCH.' / T40,
  'REDUCE STEP SIZE (CFSC) TO HAVE ONE CONSTRAINT IN CONSTRAINED SEA
  RCH.')
  GO TO 220

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C**** THIS IS POINT H WHICH APPEARS IN FLOWCHART, FIG. 13.
146 ZTEMP = ZBASE
DO 147 I = 1,N
147 XTEMP(I) = XBASE(I)
IF (KC.NE. 1) GO TO 151
148 IF(LPRINT.LE.2) WRITE (6, 1148)
1148 FORMAT(T40, 'SINGLE CONSTRAINT CONTACTED IN LOCAL SEARCH.')
```

C**** THIS IS POINT 210 WHICH APPEARS IN FLOWCHART, FIG. 14.

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210 IF (ISTOP.EQ. 2) GO TO 409
C**** THIS IS POINT M WHICH APPEARS IN FLOWCHART, FIG. 14.
220 CFSC = CFSC/3.0
IF (CFSC.GE. XLLIM) GO TO 231
IF(LPRINT.LE.3) WRITE (6, 2232)
2232 FORMAT(T40, 'INCREMENTAL VALUE OF VARIABLES .LT. LOWER LIMIT.')
```

1/T40, 'PATTERN SEARCH TERMINATED.')

```

CALL SUMS (XTEMP, ZTEMP)
IF(LPRINT.LE.3) CALL OUTPUT (XTEMP, ZTEMP)
GO TO 9999
```

C**** THIS IS POINT J WHICH APPEARS IN FLOWCHART, FIG. 13.

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149 IF(LPRINT.LE.1) WRITE (6, 1149)
1149 FORMAT(T40, 'BETTER POINT NOT FOUND IN LOCAL SEARCH.')
```

1/T40, 'REDUCE STEP SIZE (CFSC).')

C**** THIS IS POINT N WHICH APPEARS IN FLOWCHART, FIG. 13.

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215 ZTEMP = ZBASE
DO 216 I = 1,N
216 XTEMP(I) = XBASE(I)
GO TO 220
231 IF(LPRINT.LE.1) WRITE (6, 2231)
2231 FORMAT(T40, 'START LOCAL SEARCH WITH REDUCED STEP SIZE (CFSC).')
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GO TO 111

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C*** THIS IS POINT 310 WHICH APPEARS IN FLOWCHART, FIG. 17.
310 CONTINUE
329 DO 330 I = 1, N
330 XTEST(I) = XBASE(I) + ADAF*CASS*PARTH(I)
    CALL SUMS (XTEST, ZTEST)
    IF(LPRINT.LE.0) CALL OUTPUT (XTEST, ZTEST)
350 GO TO (355, 380), IFEAS
355 IF(LPRINT.LE.1) WRITE (6, 3355)
3355 FORMAT(T40, 'NO CONSTRAINT CONTACTED IN CONSTRAINED MOVE.')
360 IF (ZTEST .GT. ZBASE) GO TO 370
    IF (ADAF .LE. 1.1) GO TO 362
    IF(LPRINT.LE.1) WRITE (6, 3365)
3365 FORMAT(T40, 'ACCELERATED CONSTRAINED MOVE NOT A BETTER POINT.*/
    T40, 'LET CONSTRAINED SEARCH ACCELERATION FACTOR(ADAF) = 1.0')
    ADAF = 1.0
    GO TO 329
362 IF(LPRINT.LE.1) WRITE (6, 3362)
3362 FORMAT(T40, 'CONSTRAINT VIOLATED IN CONSTRAINED MOVE.')
    1, REDUCE STEP SIZE (CASS).')
    GO TO 381
380 IF(LPRINT.LE.1) WRITE (6, 3380)
3380 FORMAT(T40, 'CONSTRAINT CONTACTED IN CONSTRAINED SEARCH. REDUCE ST
    IEP SIZE (CASS).')
    ADAF = 1.0
381 CASS = CASS/3.0
385 IF (CASS .GT. XLLIM) GO TO 310
    IF(LPRINT.LE.2) WRITE (6, 3385)
3385 FORMAT(T40, 'CASS .LT. MIN DELTA X. CONSTRAINED SEARCH HAS FAILED.'
    1)
390 ISTOP = 1

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GO TO 215
370 ADAF = ADAF + 1.0
    ZBASE = ZTEST
    ZTEMP = ZTEST
    DO 375 I = 1, N
        X(I) = XTEST(I)
        XBASE(I) = XTEST(I)
        XTEMP(I) = XBASE(I)
375 IF(LPRINT.LE.1) WRITE (6, 3377)
3377 FORMAT(T43, 'CONSTRAINED MOVE SUCCESSFUL. START ANOTHER LOCAL SEARCH
      1H.')
```

GO TO 111

C*** THIS IS POINT P WHICH APPEARS IN FLOWCHART, FIG. 15.

```

409 DO 410 I = 1, N
410 X(I) = XBASE(I)
    CASS = CFSC
    DO 426 I = 1, N
        X(I) = X(I) + D
        CALL VAL(X, ZP)
        CP = C(LACT)
        X(I) = X(I) - 2.0 * D
        CALL VAL(X, ZM)
        CM = C(LACT)
422 PARTZ(I) = (ZP-ZM)
        PARTH(I) = (CP-CM)
        X(I) = XBASE(I)
426 CONTINUE
    DENOM = 0.0
    DO 430 I = 1, N
430 DENOM = DENOM + PARTZ(I)**2
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00780 DENOM = DSQRT(DENOM)
00781 IF (DENOM .LT. XLLIM) DENOM = XLLIM
00782 DO 440 I = 1, N
00783 PARTZ(I) = PARTZ(I)/DENOM
00784 CONTINUE
00785 IF (LPRINT.LE.2.AND.N.LE.6) WRITE(6,4441) (PARTZ(I), I=1, N)
00786 FORMAT(T40, 'OBJECTIVE GRADIENT IS', 6F10.6)
00787 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9441) (PARTZ(I), I=1, N)
00788 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00789 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00790 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00791 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00792 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00793 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00794 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00795 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00796 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00797 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00798 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00799 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00800 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00801 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00802 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00803 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00804 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00805 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00806 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00807 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00808 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)
00809 IF (LPRINT.LE.2.AND.N.GT.6) WRITE(6,9442) (PARTZ(I), I=1, N)

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IF(LPRINT.LE.2.AND.N.GT.6) WRITE(6,9495)(PARTH(I),I=1,N)
9495 FORMAT(T20,'COMBINED GRADIENT = ',9F10.6/T20,11F10.6)
IF (DENOM .GT. 0.03) GO TO 310
IF(LPRINT.LE.2) WRITE (6, 4498)
4498 FORMAT (T40,'OBJECTIVE AND CONSTRAINT NEARLY PARALLEL.')
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      ISTOP = 1
      GO TO 111
9999 CONTINUE
      RETURN
      END
      SUBROUTINE SUMS(X, Z)
      THIS SUBROUTINE APPEARS IN FLOWCHART, FIG. 4.
      IMPLICIT REAL*8(A-H,O-Z)
      DIMENSION X(20)
      COMMON C(41),CFSC,CASS,ADAF,XLLIM,KC,IFEAS,ISTOP,LACT,M,N,LPRINT
      CALL VAL(X, Z)
      KC = 0
      IFEAS = 1
      IF (M.LT. 1) GO TO 11
      DO 10 I = 1, M
      IF (C(I) .GE. 0.0) GO TO 10
      40 IFEAS = 2
      LACT = I
      50 KC = KC + 1
      10 CONTINUE
      11 CONTINUE
      RETURN
      END
      SUBROUTINE SUMC(X, CON)
      THIS SUBROUTINE APPEARS IN FLOWCHART, FIG. 5.
      C**** THIS SUBROUTINE APPEARS IN FLOWCHART, FIG. 5.

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00840      IMPLICIT REAL*8(A-H,O-Z)
00841      DIMENSION X(20)
00842      COMMON C(41),CFSC,CASS,ADAF,XLLIM,KC,IFEAS,ISTOP,LACT,M,N,LPRINT
00843      CALL VAL(X,Z)
00844      CON = 0.0
00845      IFEAS = 1
00846      DO 10 I = 1,M
00847      IF(C(I).GE. 0.0) GO TO 9
00848      IFEAS = 2
00849      CON = CON + C(I)
00850
00851      9 CONTINUE
00852      10 CONTINUE
00853      11 CONTINUE
00854      RETURN
00855      END
00856      SUBROUTINE OUTPUT(X,Z)
00857
00858      C..... THIS SUBROUTINE PRINTS OUT VALUES OF THE PATTERN SEARCH
00859      C..... VARIABLES AND CONSTRAINTS AT THE END OF EACH SUCCESSFUL STEP
00860      C..... IF LPRINT = 0.
00861      C.....
00862      IMPLICIT REAL*8(A-H,O-Z)
00863      DIMENSION X(20)
00864      COMMON C(41),CFSC,CASS,ADAF,XLLIM,KC,IFEAS,ISTOP,LACT,M,N,LPRINT
00865      IF(N.GT.6) GO TO 40
00866      WRITE (6,100) Z,(X(I),I=1,N)
00867      IF(M.GT.8) GO TO 50
00868      WRITE(6,200) (C(J),J=1,M)
00869      WRITE (6,300)
      RETURN

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50 WRITE(6,600) (C(J),J = 1,M)
   WRITE(6,300)
   RETURN
40 WRITE(6,400) Z,(X(I),I = 1,N)
   WRITE(6,500) (C(J),J = 1,M)
   WRITE(6,300)
100 FORMAT(T40,'OBJ.FCN.=',1PE12.5,' VARIAB.=',0P6F10.6)
200 FORMAT(T40,'CONSTR.=',8F10.6)
300 FORMAT(T40,60(' '))
400 FORMAT(T15,'OBJ.FCN.=',1PE12.5,' VAR.=',9F10.6,/T15,11F10.6)
500 FORMAT(T15,'CONSTR.=',10F10.6/,2(T15,10F10.6/),T15,11F10.6)
600 FORMAT(T40,'CONSTR.=',8F10.6/,4(T40,9F10.6/))
   RETURN
END
SUBROUTINE VAL (X,DUAL)
C.....
C..... THIS SUBROUTINE CALCULATES THE VALUES OF ALL THE DUAL VARIABLES,
C..... GIVEN THE LAST IDDIFF OF THEM, AND THEN CALCULATES THE VALUE
C..... OF THE DUAL FUNCTION.
C.....
C..... IMPLICIT REAL*8(A-B,D-H,O-Q,U-Z), INTEGER(C,R-T)
      DIMENSION X(20)
      COMMON/TOY/A(10,10,10),B(31,32),XCON(10,10),SIGMA(10,10),SIG(10),
      1RJCW(32),CJL(32),TSJBM(10),CAPM,CAPN,CAPNP1,CAPNP2,CAPT,CAPTPI
      COMMON W(41),XCFSC,XCASS,ADAF,XLLIM,KC,IFEAS,ISTOP,LACT,MPT,IDDIF,
      1LPRINT
      IF(IDDIFF.EQ.0) X(1) = 0.0
      DO 1 I=1,CAPNP1
        BSUM = B(I,CAPTPI)
        DO 2 J=CAPNP2,CAPT

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C.....THE LAST IDDIFF DUAL VARIABLES W = THE PATTERN SEARCH VARIABLES X.
C.....
C.....
C.....W(J)=X(J-CAPNP1)
C.....2 BSUM = BSUM - B(I,J)*W(J)
C.....
C.....SOLVE FOR THE FIRST DUAL VARIABLES.
C.....
C.....1 W(I) = BSUM
C.....IF(CAPM.EQ.1) GO TO 6
C.....I = TSUBM(1)
C.....
C.....THERE IS ONE ADDITIONAL DUAL CONSTRAINT FOR EACH PRIMAL
C.....CONSTRAINT.
C.....
C.....DO 3 IM = 2,CAPM
C.....DSUM = 0.0
C.....ITS = TSUBM(IM)
C.....DO 4 IT = 1,ITS
C.....I=I+1
C.....DS = SIGMA(IM,IT)
C.....4 DSUM = DSJM+DS*W(I)
C.....DS = SIG(IM)
C.....3 W(CAPT+IM) = DSUM*DS
C.....6 W(CAPT+1) = 1.0
C.....
C.....SOLVE FOR THE DUAL FUNCTION.
C.....
C.....I = 0
C.....DUAL = 1.0

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00930      DO 5 IM = 1,CAPM
00931      ITS = TSBM(IM)
00932      DO 5 IT = 1,ITS
00933      I = I+1
00934      DTERM = 1.0
00935
00936      C.....
00937      C..... IF W(I).= 0.0, THE TERM EQUALS ONE.
00938      C.....
00939      IF(W(I)) 10,5,10
00940      10 Y = XCON(IM,IT)*W(CAPT+IM)/W(I)
00941      IF(Y) 5,11,12
00942      11 DTERM = 0.0
00943      IF(SIG(1).LE.0) DTERM = 10.**(-30)
00944      GO TO 5
00945      12 Z = SIGMA(IM,IT)*W(I)
00946      DTERM = Y**Z
00947      5 DUAL = DUAL*DTERM
00948      IF(DUAL) 20,21,22
00949      20 DUAL = -1.0*(10.**(-30))
00950      RETURN
00951      21 DUAL = 0.0
00952      IF(SIG(1).LE.0) DUAL = 10.**(-30)
00953      RETURN
00954      22 DUAL = SIG(1)*(DUAL**SIG(1))
00955      RETURN
00956      END
00957      BLOCK DATA
00958      IMPLICIT REAL*8(A-B,D-H,O-Q,U-Z), INTEGER(C,R-T)
00959      COMMON/TOM/A(10,10,10),B(31,32),XCON(10,10),SIGMA(10,10),SIG(10),
      1ROW(32),CJL(32),TSBM(10),CAPM,CAPN,CAPNP1,CAPNP2,CAPT,CAPTPI

```



```
COMMON/GUS/X(20),XBASE(20),XTEMP(20),XTEST(20),PARTZ(20),PARTH(20)
DATA A,XCJN,X,SIGMA,ROW,COL/1100*0.0,20*0.2,100*0,
1 1,2,3,4,5,6,7,8,9,10,11,12,13,14,15,16,17,18,19,20,21,22,23,24,25
1,26,27,28,29,30,31,32,
1 1,2,3,4,5,6,7,8,9,10,11,12,13,14,15,16,17,18,19,20,21,22,23,24,25
1,26,27,28,29,30,31,32/
END
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APPENDIX B
EXAMPLE OF OUTPUT

POLYNOMIAL MINIMIZATION BY GEOMETRIC PROGRAMMING

SEA POWER PROBLEM - CUFFIN, ZENER, AND PETERSON - 10^{**6} WATTS

PRIMAL PROGRAM.

THE NUMBER OF PRIMAL VARIABLES = 7
THE NUMBER OF PRIMAL CONSTRAINTS = 4

*** GENERALIZED POLYNOMIAL ***

ROW	TERM	COEFFICIENT OF TERM	PRIMAL VARIABLE	EXPONENT OF VARIABLE	SIGNUM OF ROW
1	1	1.0000000	* X(1)	**	1.0000000

MINIMIZE.

1 1 1.0000000 * X(1) ** 1.0000000

SIGN OF PRIMAL FUNCTION AT MINIMUM ASSUMED TO BE POSITIVE

SUBJECT TO

2 1 1.0000000 * X(4) ** 1.0000000

.PLUS.

2 2 1.0000000 * X(5) ** 1.0000000

.PLUS.

2 3 1.0000000 * X(6) ** 1.0000000

.PLUS.

2 4 1.0000000 * X(7) ** 1.0000000

.LE. 1

3 1 0.180000000 * X(1) ** -1.0000000
 * X(2) ** 1.0000000
 * X(6) ** -1.0000000

.LE. 1

4 1 40.0000000 * X(1) ** -1.0000000
 * X(2) ** 1.0000000
 * X(3) ** -1.0000000
 * X(5) ** -1.0000000

.LE. 1

5 1 445000000. * X(2) **
* X(4) ** -1.0000000
-1.0000000

.PLUS.

5 2 0.600000000D-07 * X(1) **
* X(2) ** 1.0000000
* X(3) ** -1.0000000
* X(4) ** 3.0000000
-1.0000000

.PLUS.

5 3 0.215000000D-07 * X(3) **
* X(4) ** 2.0000000
* X(7) ** -1.0000000
-1.0000000

.LE. 1

THE TOTAL NUMBER OF TERMS = 10
THE NUMBER OF 'DEGREES OF DIFFICULTY' = 2

DUAL PROGRAM.

THE NUMBER OF DUAL VARIABLES = 10
THE NUMBER OF DUAL EQUATIONS = 8

*** EXTENDED MATRIX OF COEFFICIENTS OF DUAL VARIABLES ***

1.000	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	1.000
0.0	1.000	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	-1.000	-1.000	-1.000	0.0	0.0	1.000
0.0	0.0	1.000	0.0	0.0	0.0	0.0	0.0	0.0	0.0	0.0	-3.000	-2.000	-2.000	0.0	0.0	0.0
0.0	0.0	0.0	1.000	0.0	0.0	0.0	0.0	0.0	0.0	0.0	2.000	2.000	2.000	0.0	0.0	1.000
0.0	0.0	0.0	0.0	1.000	0.0	0.0	0.0	0.0	0.0	0.0	0.0	-1.000	-1.000	0.0	0.0	0.0
0.0	0.0	0.0	0.0	0.0	1.000	0.0	0.0	0.0	0.0	0.0	2.000	2.000	2.000	0.0	0.0	1.000
0.0	0.0	0.0	0.0	0.0	0.0	1.000	0.0	0.0	0.0	0.0	-3.000	-2.000	-2.000	0.0	0.0	0.0
0.0	0.0	0.0	0.0	0.0	0.0	0.0	1.000	0.0	0.0	0.0	0.0	0.0	0.0	1.000	0.0	1.000

PATTERN SEARCH.

[illegible]

```
THE NUMBER OF VARIABLES = 2
THE NUMBER OF CONSTRAINTS = 14
THE LOWER LIMIT ON STEP SIZE = 1.0000D-06
PRINT OPTION = 2
```

INITIAL POINT.

OBJ.FCN. = 8.60619D 07	VARIAB. =	0.200000	0.200000
CONSTR. =	1.000000	1.000000	0.200000
0.200000	0.200000	1.000000	0.200000
		2.800000	1.400000
		0.200000	1.000000
		0.200000	0.200000
		1.000000	0.200000
		1.000000	1.000000

INCREMENTAL VALUE OF VARIABLES .LT. LOWER LIMIT.

PATTERN SEARCH TERMINATED.

OBJ.FCN. =	1.267480	08	VARIAB. =	0.246988	0.042763	
CONSTR. =	1.000000	1.289751		0.826490	0.420498	0.042763
	0.246988	0.042763	1.000000	2.579502	0.420498	0.826490
						1.289751

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DUAL SOLUTION.

THE VALUE OF THE DUAL FUNCTION AT A MAXIMUM POINT =
0.12674758D 09

THE CORRESPONDING DUAL VARIABLES ARE.

DUAL VARIABLE	VALUE
w(1) =	1.0000000
w(2) =	1.2897512
w(3) =	0.82649018
w(4) =	0.42049763
w(5) =	0.42763375D-01
w(6) =	0.42049763
w(7) =	0.82649018
w(8) =	1.0000000
w(9) =	0.24698781
w(10) =	0.42763375D-01
w(11) =	1.0000000
w(12) =	2.5795024
w(13) =	0.42049763
w(14) =	0.82649018
w(15) =	1.2897512

PRIMAL SOLUTION.

THE VALUES OF THE PRIMAL VARIABLES AT A MINIMUM POINT ARE -----

PRIMAL VARIABLE	VALUE
X(1) =	0.12674758D 09
X(2) =	0.11478755D 09
X(3) =	113.06113
X(4) =	0.50000000
X(5) =	0.32040683
X(6) =	0.16301502
X(7) =	0.16578149D-01

THE CONTRIBUTION OF EACH TERM TO THE PRIMAL FUNCTION IS.

TERM CONTRIBUTION

1 0.12674758D 09

THE VALUE OF ROW 1 = 0.12674758D 09

2 0.50000000

3 0.32040683

4 0.16301502

5 0.16578149D-01

THE VALUE OF ROW 2 = 1.00000000

6 1.00000000

THE VALUE OF ROW 3 = 1.00000000

7 1.00000000

THE VALUE OF ROW 4 = 1.00000000

8 0.77534539

9 0.19149884

10 0.33155768D-01

THE VALUE OF ROW 5 = 1.00000000

THE VALUE OF THE PRIMAL FUNCTION = 0.12674758D 09
THE VALUE OF THE DUAL FUNCTION = 0.12674758D 09
THE ABSOLUTE VALUE OF THE DIFFERENCE = 0.10952353D-05
THE RELATIVE DIFFERENCE = 0.109524D-05/ 0.126748D 09
= 0.000000 PER CENT

PRIMAL PROBLEM IS CONVEX. GLOBAL MINIMUM GUARANTEED.

ELAPSED TIME = 1.810 SECONDS.